
A combinatorial proof of Macdonald positivity

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Macdonald positivity

In 1988, Macdonald defined $\tilde{H}_\mu(x; q, t)$ and proved it to be symmetric. Writing $\tilde{H}_\mu(x; q, t) = \sum_\lambda \tilde{K}_{\lambda, \mu}(q, t) s_\lambda(x)$, he conjectured:

Theorem. (Haiman 2001) *We have $\tilde{K}_{\lambda, \mu}(q, t) \in \mathbb{N}[q, t]$.*

Original proof realizes $\tilde{H}_\mu(x; q, t)$ as the Frobenius series of doubly graded \mathcal{S}_n -module (Garsia-Haiman module) using algebraic geometry of the Hilbert scheme.

Newer proof due to Grojnowski and Haiman relates $\tilde{H}_\mu(x; q, t)$ to LLT positivity (Haglund's formula) and then uses Kazhdan-Lusztig theory.

Problem: Find a *combinatorial* proof of positivity.

A monomial expansion for $\tilde{H}_\mu(x; q, t)$

Theorem. (Haglund, Haiman, Loehr 2005)

$$\begin{aligned}\tilde{H}_\mu(x; q, t) &= \sum_{S: \mu \rightarrow \mathbb{N}} q^{\text{inv}(S)} t^{\text{maj}(S)} x^S \\ &= \sum_{S: \mu \rightsquigarrow [n]} q^{\text{inv}(S)} t^{\text{maj}(S)} Q_{\sigma(S)}(x)\end{aligned}$$

Proposition. (Gessel 1984)

$$\begin{aligned}s_\lambda(x) &= \sum_{T \in \text{SSYT}(\lambda)} x^T \\ &= \sum_{T \in \text{SYT}(\lambda)} Q_{\sigma(T)}(x)\end{aligned}$$

Quasi-symmetric functions

For $\sigma \in \{\pm 1\}^{n-1}$, define the *quasi-symmetric function*

$$Q_\sigma(x) = \sum_{\substack{i_1 \leq \dots \leq i_n \\ i_j = i_{j+1} \Rightarrow \sigma_j = +1}} x_{i_1} \cdots x_{i_n}.$$

Define the *descent signature* $\sigma : \text{SYT} \rightarrow \{\pm 1\}^{n-1}$ by

$$\sigma(T)_i = \begin{cases} +1 & i \text{ left of } i+1 \text{ in } w_T \\ -1 & i+1 \text{ left of } i \text{ in } w_T \end{cases}$$

$$\sigma \left(\begin{array}{|c|c|c|} \hline 5 & 7 & 10 \\ \hline 2 & 6 & 8 \\ \hline 1 & 3 & 4 & 9 \\ \hline \end{array} \right) = - + + - + - + + -$$

Making a vertex-signed graph

Given a quasi-symmetric expansion

$$G(x) = \sum_{v \in V} Q_{\sigma(v)}(x),$$

define a vertex-signed graph \mathcal{G} with vertex set V and signature function $\sigma : V \rightarrow \{\pm 1\}^{n-1}$.

Goal: Give sufficient conditions for a vertex-signed graph $\mathcal{G} = (V, \sigma, E)$ to have connected components which satisfy

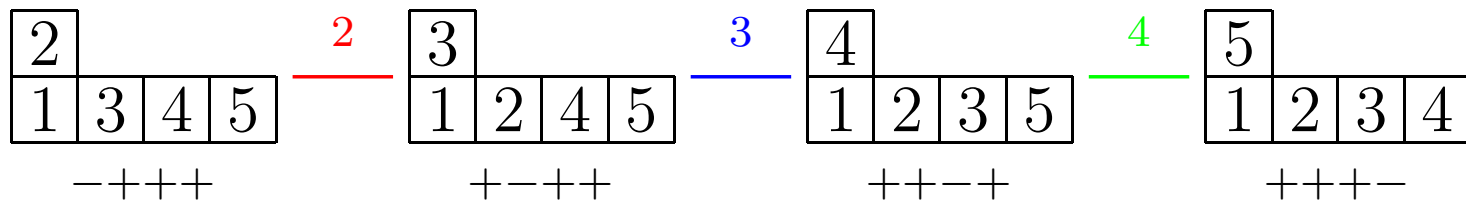
$$\sum_{v \in \mathcal{C}} Q_{\sigma(v)}(x) = s_{\lambda}(x).$$

Dual equivalence and graphs

An *elementary dual equivalence* for $i-1, i, i+1$ on a standard word is given by

$$\dots \color{red}{i} \dots \color{green}{i \pm 1} \dots \color{blue}{i \mp 1} \dots \equiv^* \dots \color{blue}{i \mp 1} \dots \color{green}{i \pm 1} \dots \color{red}{i} \dots$$

For $T, U \in \text{SYT}$, connect T and U with an i -colored edge whenever w_T and w_U differ by an elementary dual equivalence for $i-1, i, i+1$.



Standard dual equivalence graphs

Proposition. (Haiman 1992) For T, U of partition shape,

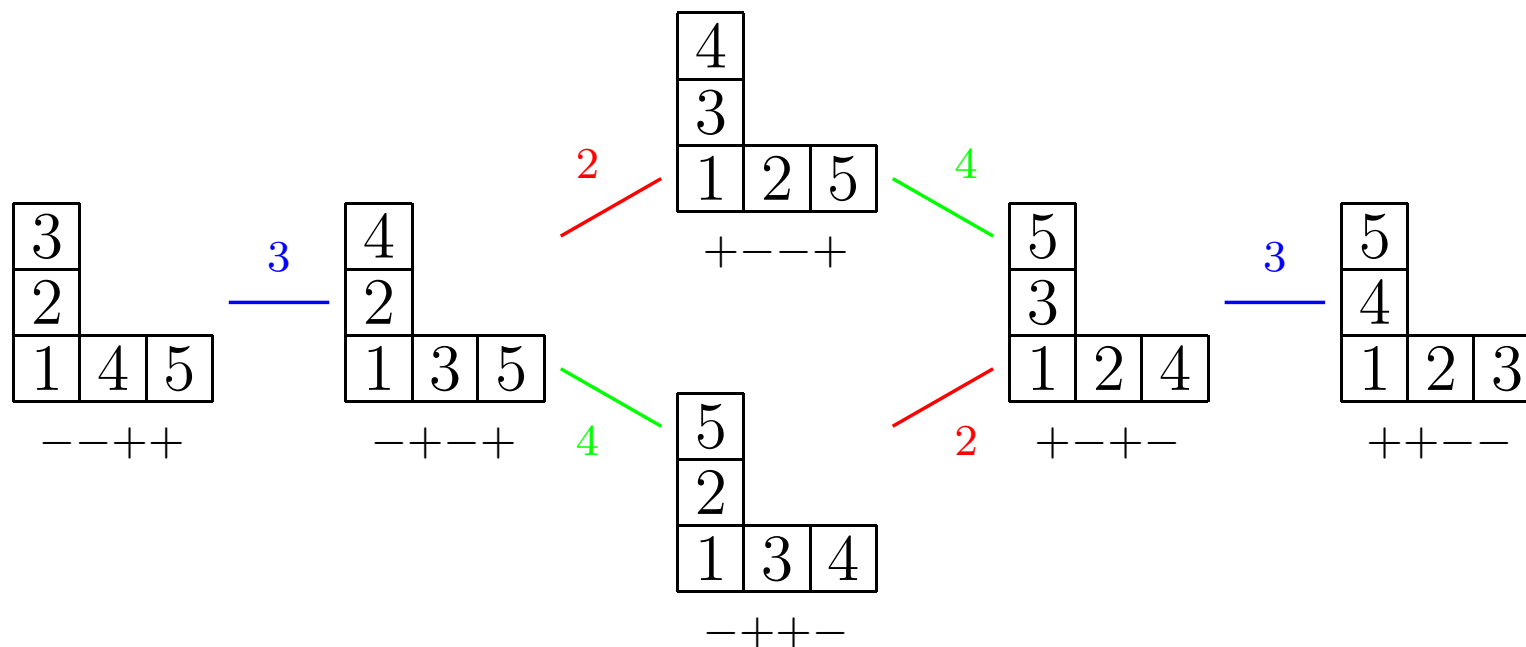
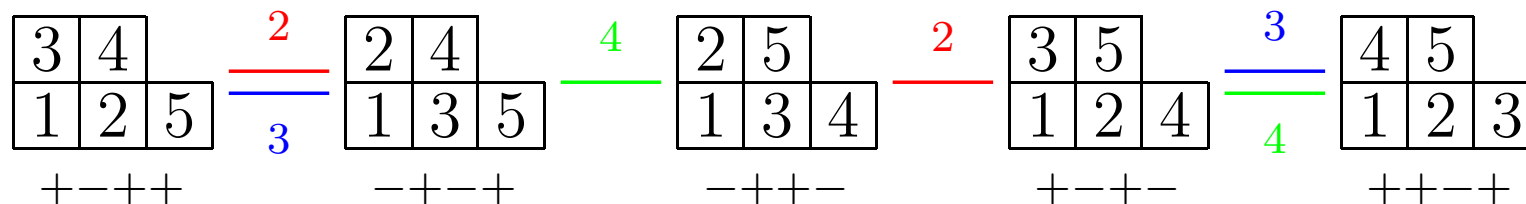
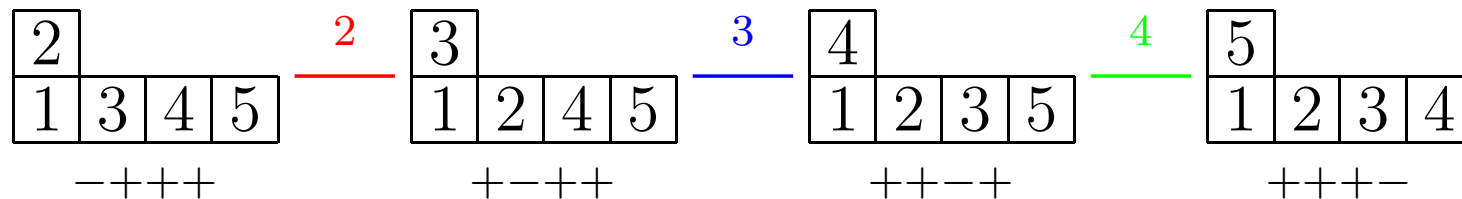
$$T \equiv^* U \Leftrightarrow \text{shape}(T) = \text{shape}(U).$$

Definition. Let \mathcal{G}_λ denote the connected component of the graph containing $\text{SYT}(\lambda)$.

The generating function of \mathcal{G}_λ is $\sum_{T \in \text{SYT}(\lambda)} Q_{\sigma(T)}(x) = s_\lambda(x)$.

Goal: Give sufficient conditions for a vertex-signed, edge-colored graph $\mathcal{G} = (V, \sigma, E)$ to have connected components isomorphic to \mathcal{G}_λ .

Examples of \mathcal{G}_λ



Dual equivalence graphs

Definition. A vertex-signed, edge-colored graph \mathcal{G} is a *dual equivalence graph* (DEG) if it satisfies 5 local axioms about signatures and edge colors.

Proposition. *The graph \mathcal{G}_λ is a DEG.*

Theorem. *Every connected component of a DEG is isomorphic to \mathcal{G}_λ for a unique partition λ .*

Corollary. *If \mathcal{G} is a DEG and α, β are statistics on $V(\mathcal{G})$ which are constant on connected components, then*

$$\sum_{v \in V(\mathcal{G})} q^{\alpha(v)} t^{\beta(v)} Q_{\sigma(v)}(x) = \sum_{\lambda} \left(\sum_{\mathcal{C} \cong \mathcal{G}_\lambda} q^{\alpha(\mathcal{C})} t^{\beta(\mathcal{C})} \right) s_\lambda(x).$$

Back to Macdonald polynomials

$$\tilde{H}_\mu(x; q, t) = \sum_{S: \mu \rightarrow [n]} q^{\text{inv}(S)} t^{\text{maj}(S)} Q_{\sigma(S)}(x)$$

$V_\mu = \{\text{standard fillings of } \mu\}$

$\sigma : V_\mu \rightarrow \{\pm 1\}^{n-1}$ using row reading word

E must preserve **inv** and **maj**

$\mathcal{H}_\mu = (V_\mu, \sigma, E)$ must be a **DEG**

The inv and maj statistics

For S a filling of the Young diagram of μ , define

$$\text{maj}(S) = |\text{Des}(S)| + \sum_{c \in \text{Des}(S)} l(c),$$

$$\text{inv}(S) = |\text{Inv}(S)| - \sum_{c \in \text{Des}(S)} a(c).$$

$$\text{Des} \left(\begin{array}{|c|c|c|} \hline 6 & 7 & 5 \\ \hline 4 & 1 & 10 \\ \hline 3 & 9 & 2 & 8 \\ \hline \end{array} \right) = \{6, 7, 4, 10\}$$

$$\text{Inv} \left(\begin{array}{|c|c|c|} \hline 6 & 7 & 5 \\ \hline 4 & 1 & 10 \\ \hline 3 & 9 & 2 & 8 \\ \hline \end{array} \right) = \left\{ \begin{array}{l} (6, 5), (7, 5), (7, 4), (5, 4), \\ (5, 1), (10, 3), (10, 9), \\ (3, 2), (9, 2), (9, 8) \end{array} \right\}$$

Edges of the graph

Define involutions

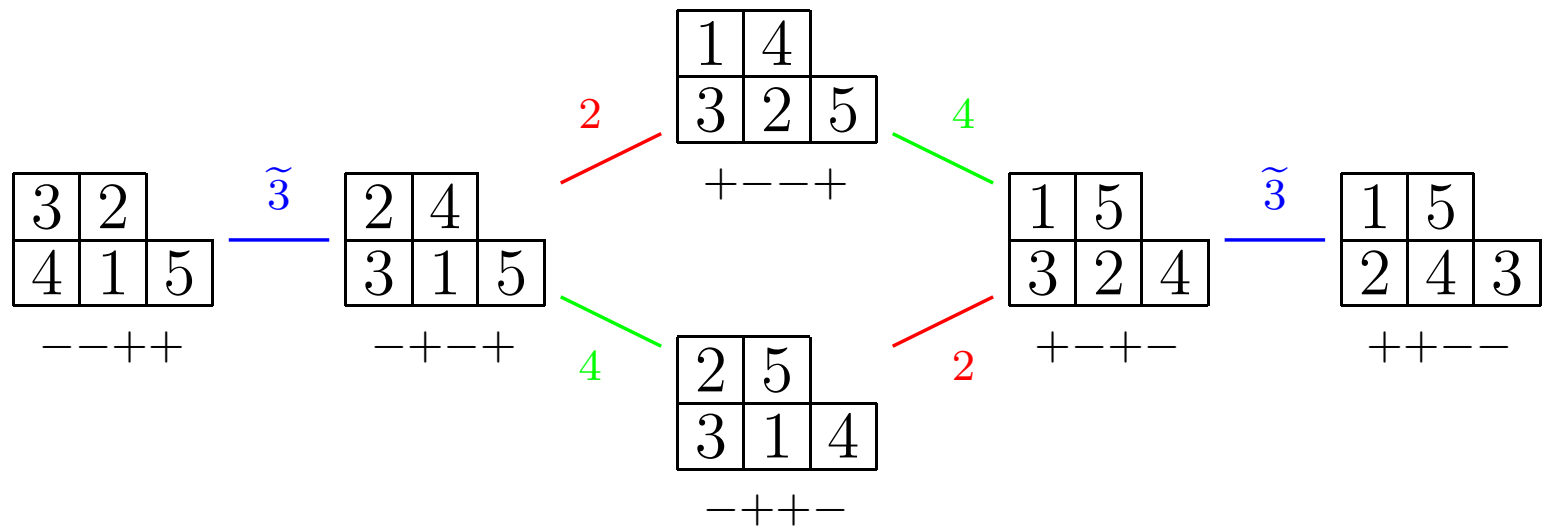
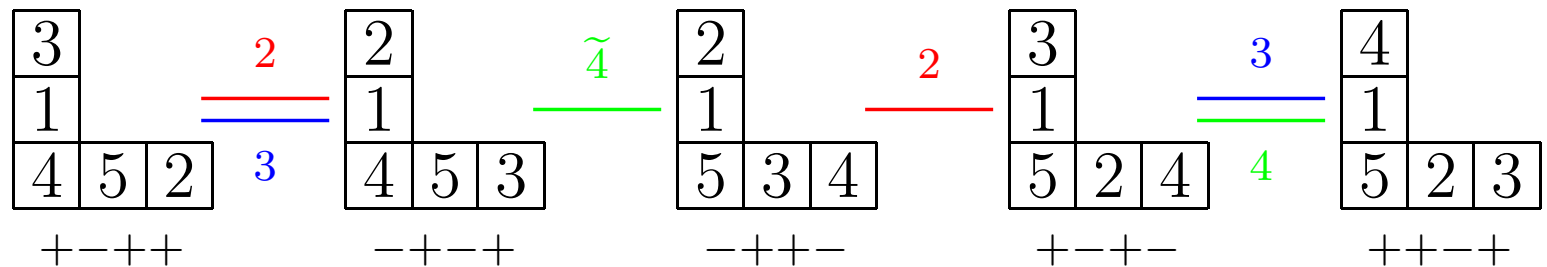
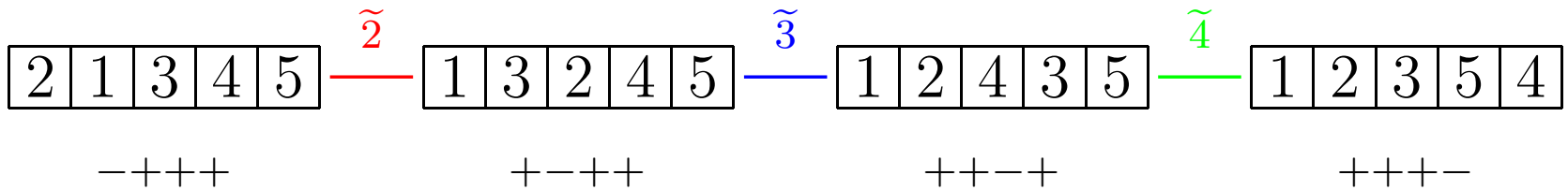
$$\begin{array}{ccccccc}
 i & i \pm 1 & i \mp 1 & \xleftrightarrow{d_i} & i \mp 1 & i \pm 1 & i \\
 i & i \pm 1 & i \mp 1 & \xleftrightarrow{\tilde{d}_i} & i \pm 1 & i \mp 1 & i
 \end{array}$$

$$D_i(S) = \begin{cases} \tilde{d}_i(S) & \text{if } i-1, i, i+1 \text{ fit in } \begin{array}{|c|c|c|c|} \hline & & & \\ \hline & \cdots & & \\ \hline \end{array} \\ d_i(S) & \text{otherwise.} \end{cases}$$

Note: $\text{Des}(S) = \text{Des}(D_i(S))$ and $|\text{Inv}(S)| = |\text{Inv}(D_i(S))|$.

Define i -colored edges on V_μ by $S \xrightarrow{i} D_i(S)$.

Examples of \mathcal{H}_μ

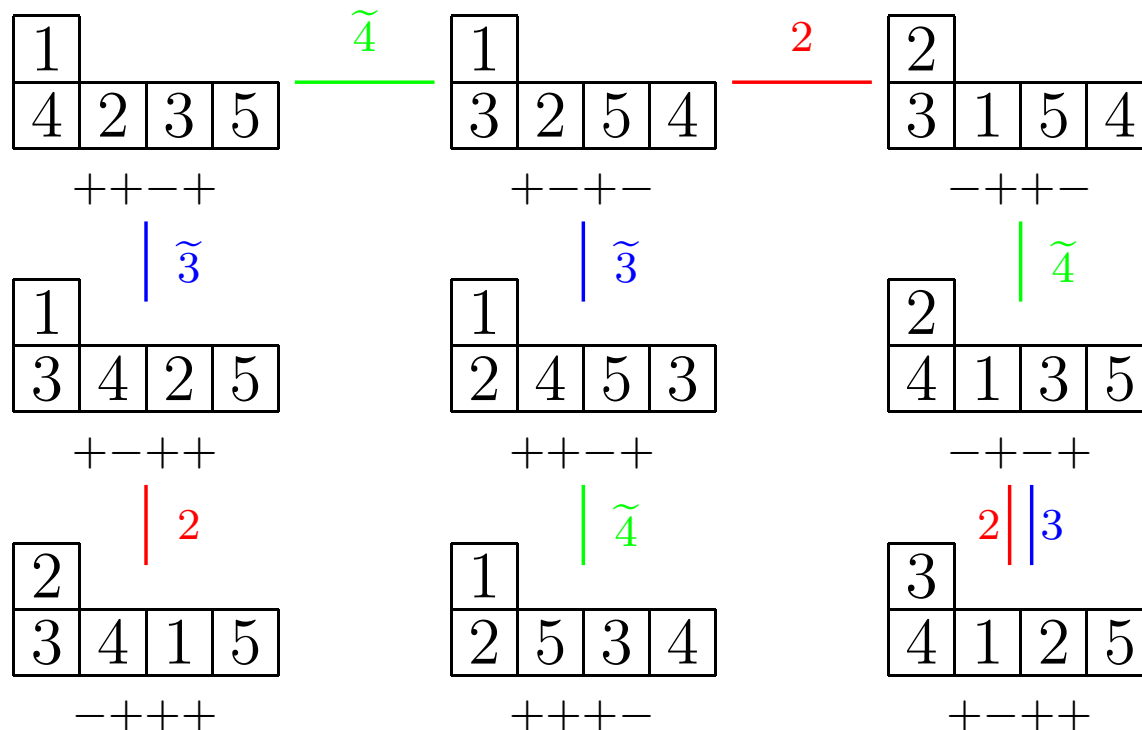


Almost a DEG

Theorem. For $\mu_1 \leq 2$, the graph \mathcal{H}_μ is a **DEG** for which **maj** and **inv** are constant on connected components.

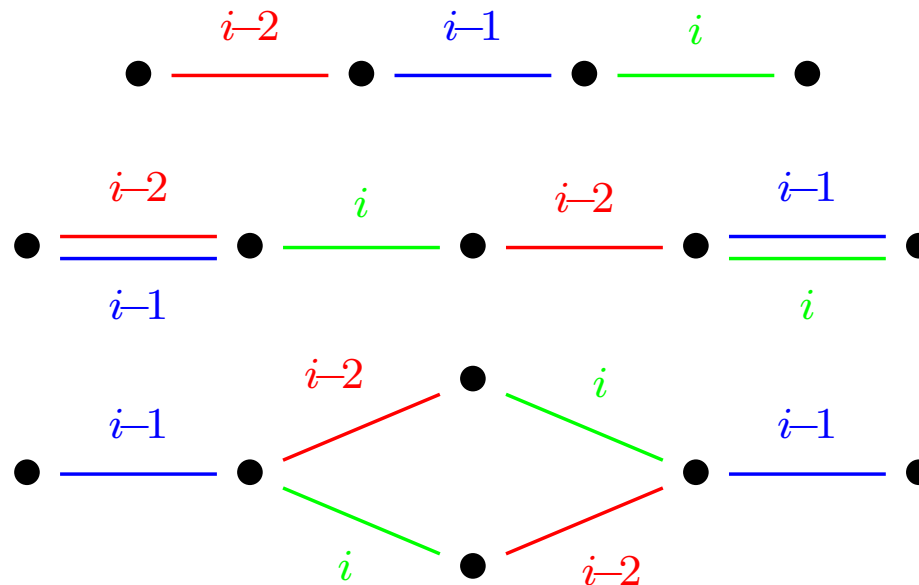
In general, \mathcal{H}_μ satisfies **DEG** axioms 1, 2, 3 and 5.

For $\mu_1 \geq 3$, \mathcal{H}_μ fails to satisfy axiom 4, so \mathcal{H}_μ is **not** a **DEG**.



Generalizing DEG axiom 4

Axiom 4: Every nontrivial connected component of $(V, \sigma, E_{i-2} \cup E_{i-1} \cup E_i)$ is one of the following:



Axiom 4': Every nontrivial connected component of $(V, \sigma, E_{i-2} \cup E_{i-1} \cup E_i)$ is in bijection with $\bigcup_{\lambda_j} \text{SYT}(\lambda_j)$ for distinct partitions λ_j of 5, and this bijection, say ϕ , satisfies $\sigma(v)_{i-3, i-2, i-1, i} = \sigma(\phi(v))_{1, 2, 3, 4}$ for all v .

Definition. A vertex-signed, edge-colored graph \mathcal{G} is a *D graph* if it satisfies axioms 1, 2, 3, 4' and 5.

Theorem. The graph \mathcal{H}_μ is a *D graph* for which **maj** and **inv** are constant on connected components.

Theorem. Let $\mathcal{G} = (V, \sigma, E)$ be a *D graph* with statistics α and β which are constant on connected components of \mathcal{G} . Then there exists a **DEG** $\tilde{\mathcal{G}} = (V, \sigma, \tilde{E})$ such that α and β are constant on connected components of $\tilde{\mathcal{G}}$.

The proof of the theorem is constructive. In fact, the transformation from \mathcal{G} to $\tilde{\mathcal{G}}$ is canonical.

Constructing a DEG from a D graph

Let \mathcal{G} be a **D graph**. We construct a sequence of graphs

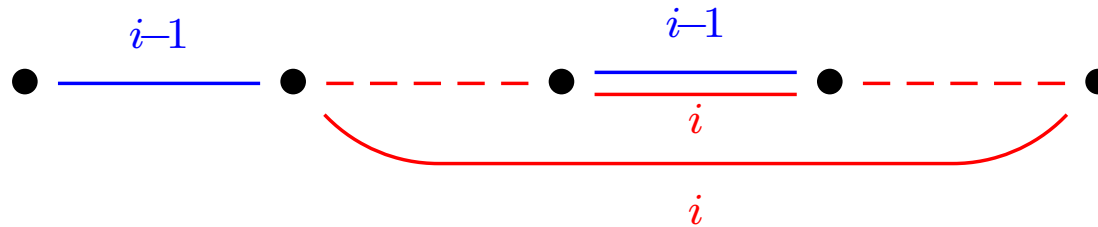
$$\mathcal{G} = \mathcal{G}_2, \mathcal{G}_3, \dots, \mathcal{G}_{n-1} = \tilde{\mathcal{G}}$$

with the following properties:

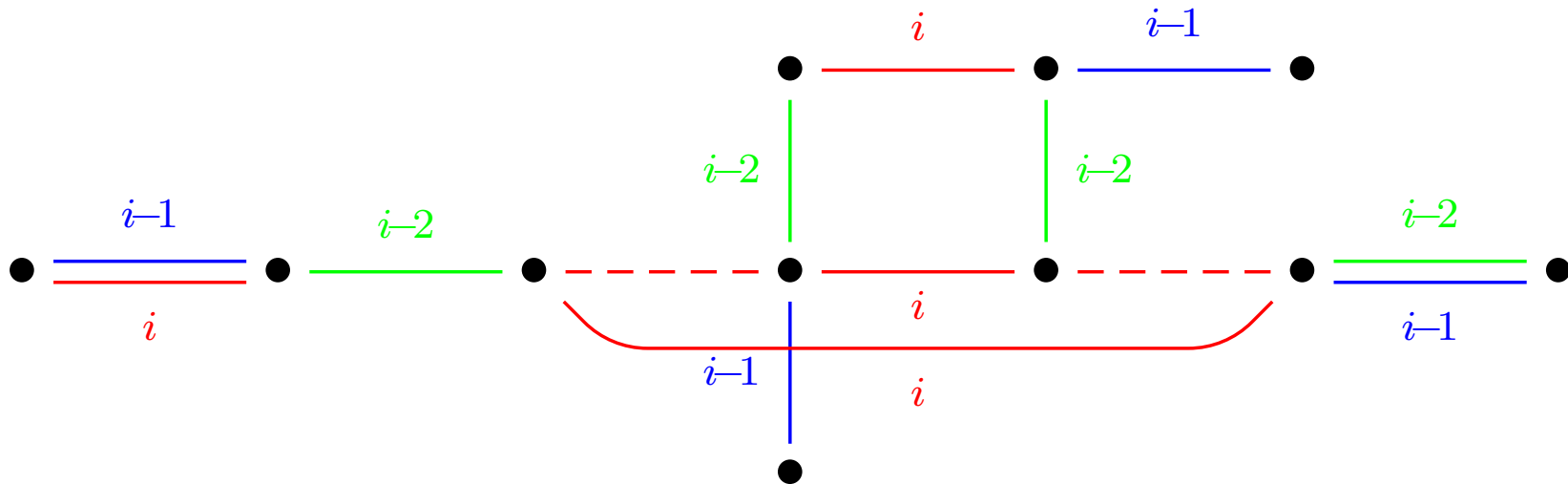
- the vertex set is the same for \mathcal{G}_i and \mathcal{G}_{i-1} ;
- the signature function is the same for \mathcal{G}_i and \mathcal{G}_{i-1} ;
- for $j \neq i$, the j -edges are the same for \mathcal{G}_i and \mathcal{G}_{i-1} ;
- if $w \overset{i}{\text{---}} x$ in \mathcal{G}_i , then w and x lie on the same connected component of $(V, \sigma, E_2 \cup \dots \cup E_i)$ in \mathcal{G}_{i-1} ;
- each \mathcal{G}_i is a **D graph**, and its restriction to $(V, \sigma, E_2 \cup \dots \cup E_i)$ is a **DEG**.

Constructing \mathcal{G}_i from \mathcal{G}_{i-1}

Step 1: Fix connected components of $(V, \sigma, E_{i-1} \cup E_i)$.

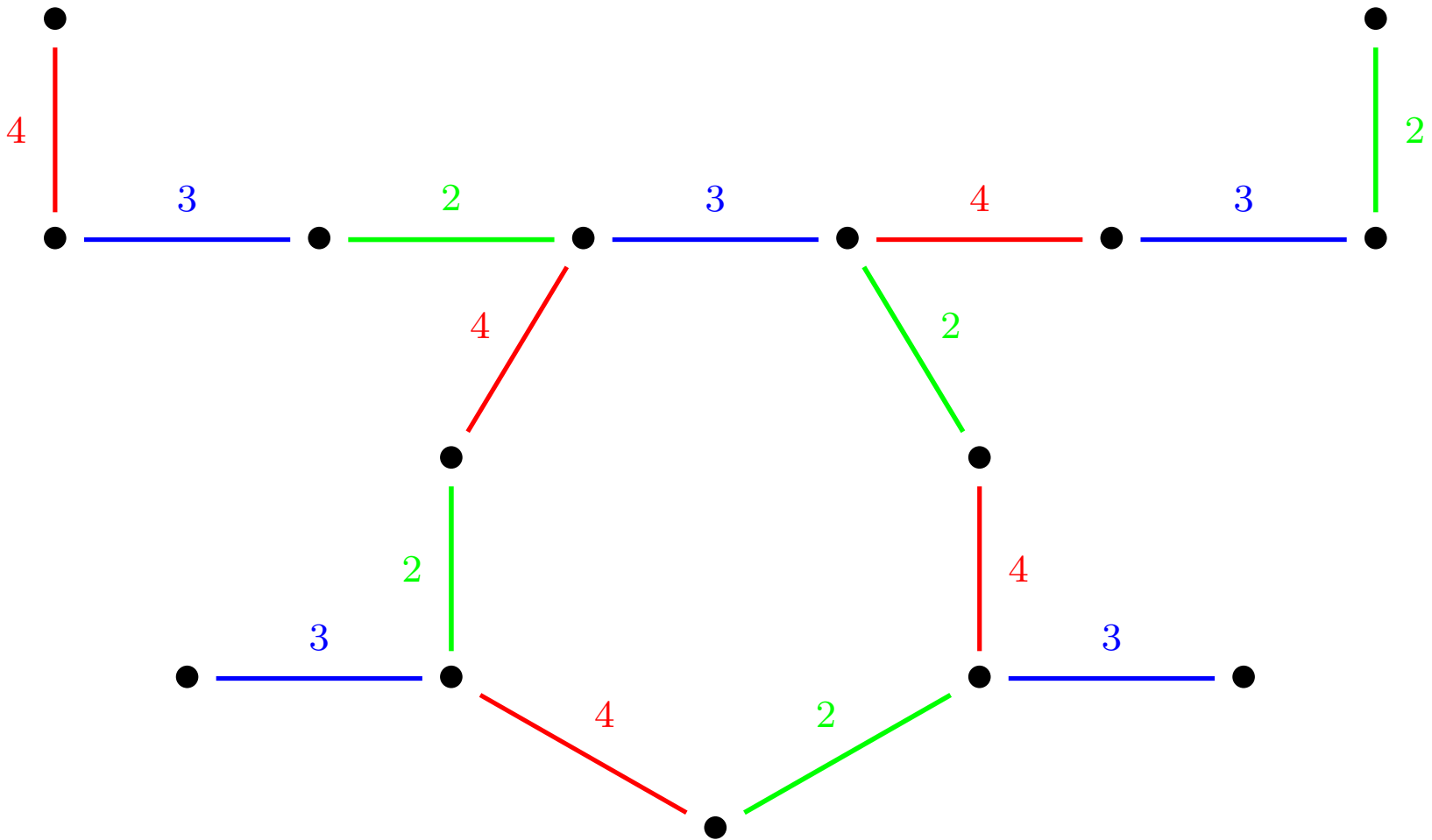


Step 2: Fix connected components of $(V, \sigma, E_{i-2} \cup E_{i-1} \cup E_i)$.



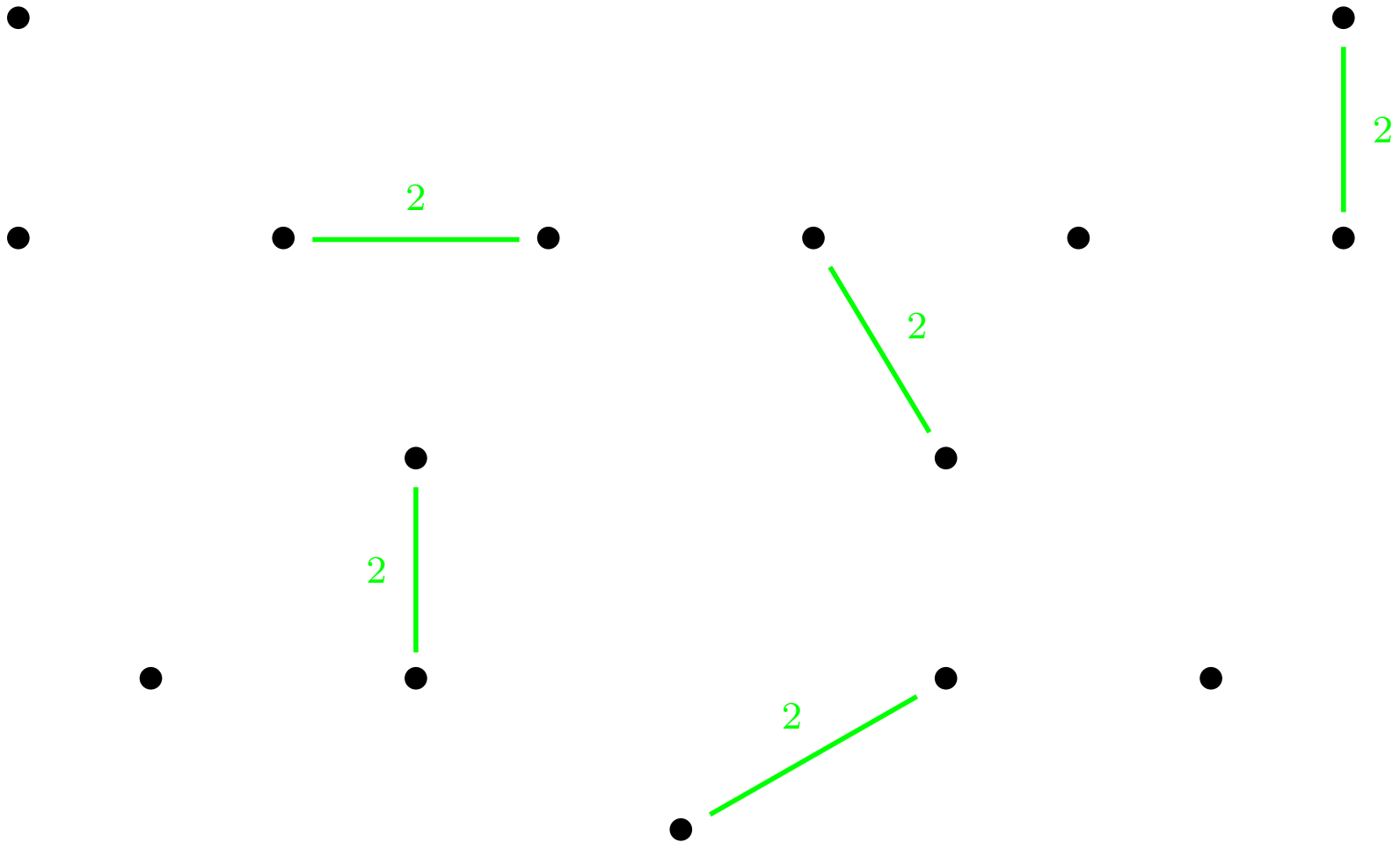
A final example: $\mathcal{G} \rightsquigarrow \tilde{\mathcal{G}}$

$$\mathcal{G} = \mathcal{G}_2 \rightsquigarrow \mathcal{G}_{2.5} \rightsquigarrow \mathcal{G}_3 \rightsquigarrow \mathcal{G}_{3.5} \rightsquigarrow \mathcal{G}_4 = \tilde{\mathcal{G}}$$



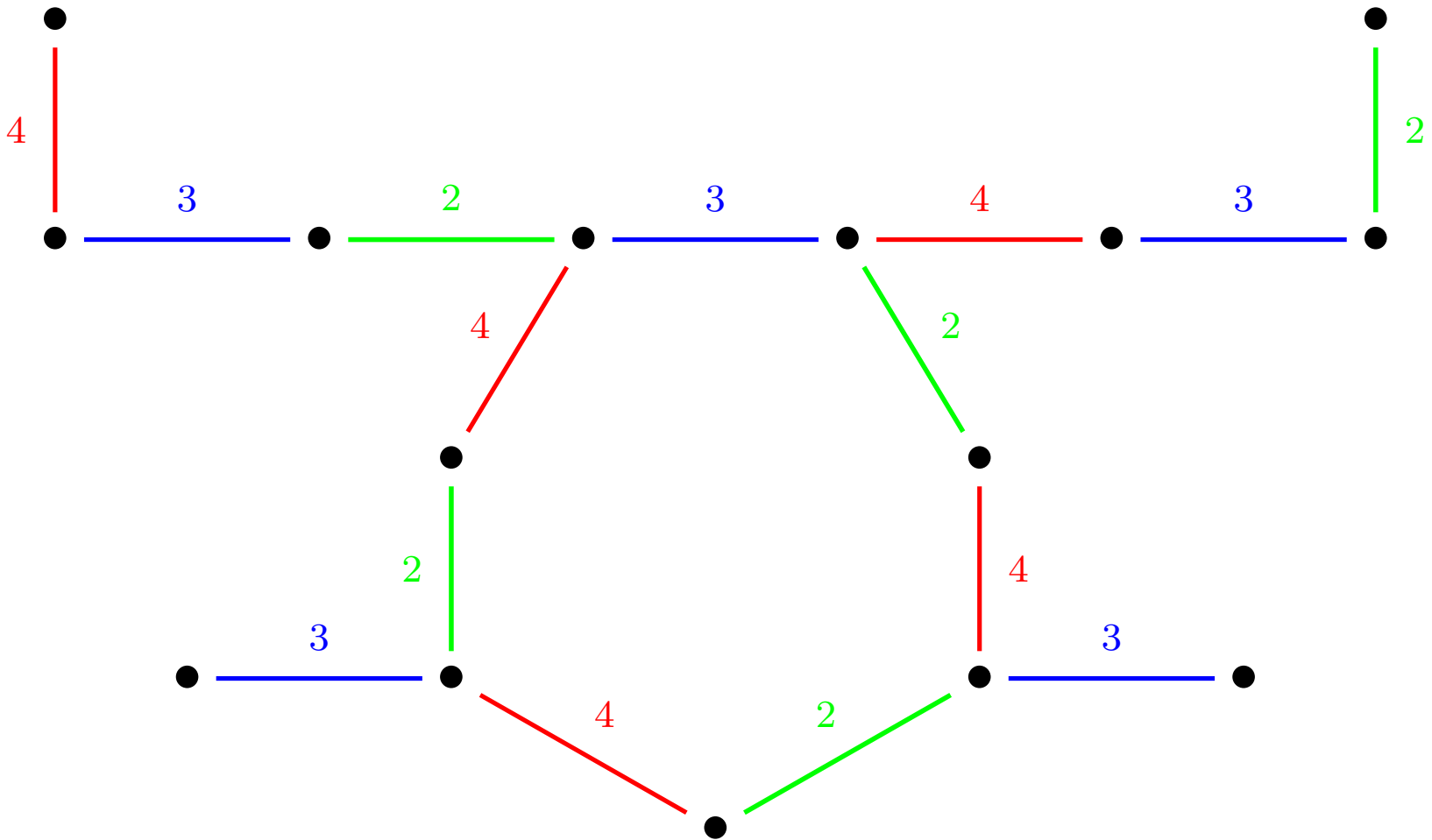
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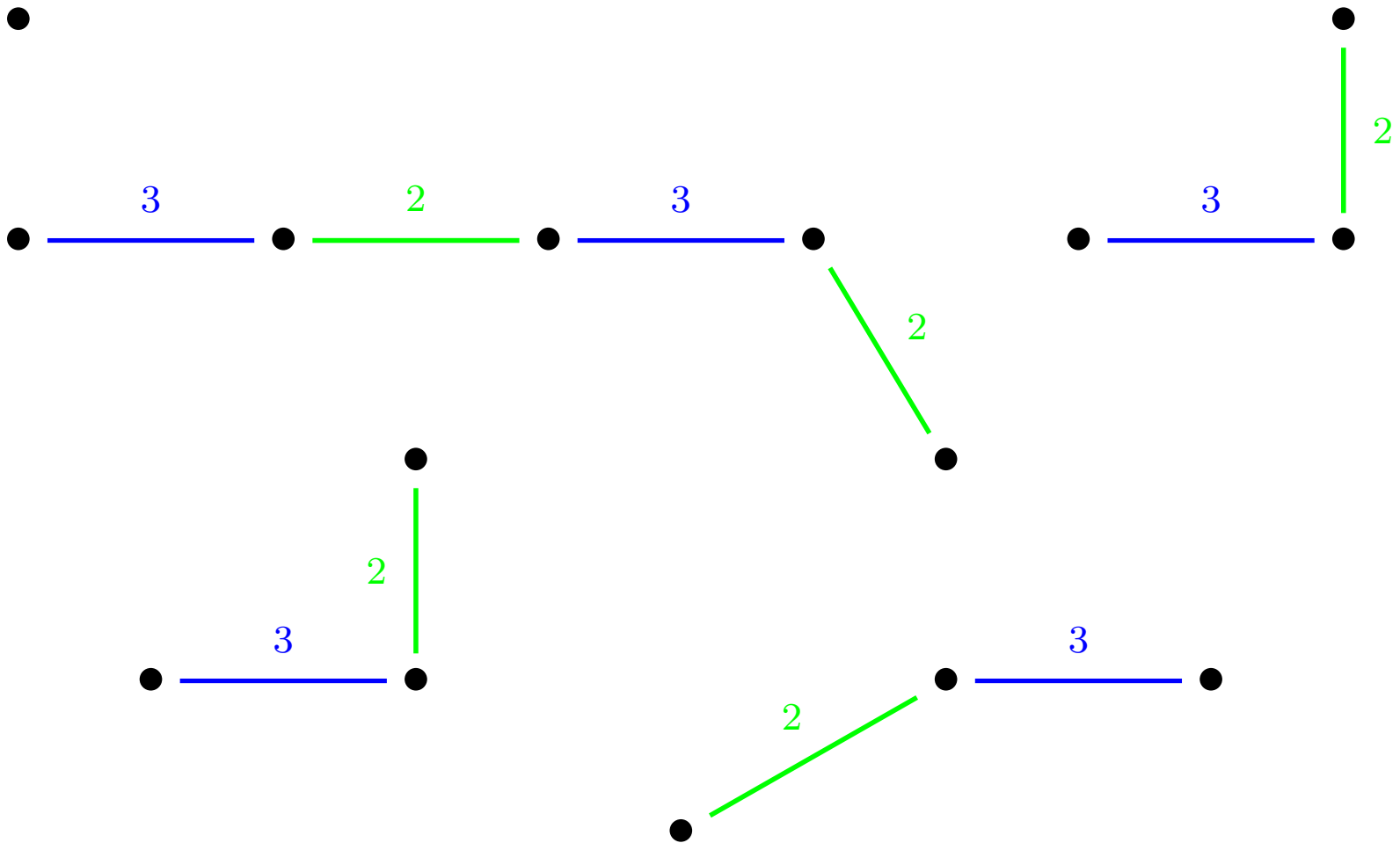
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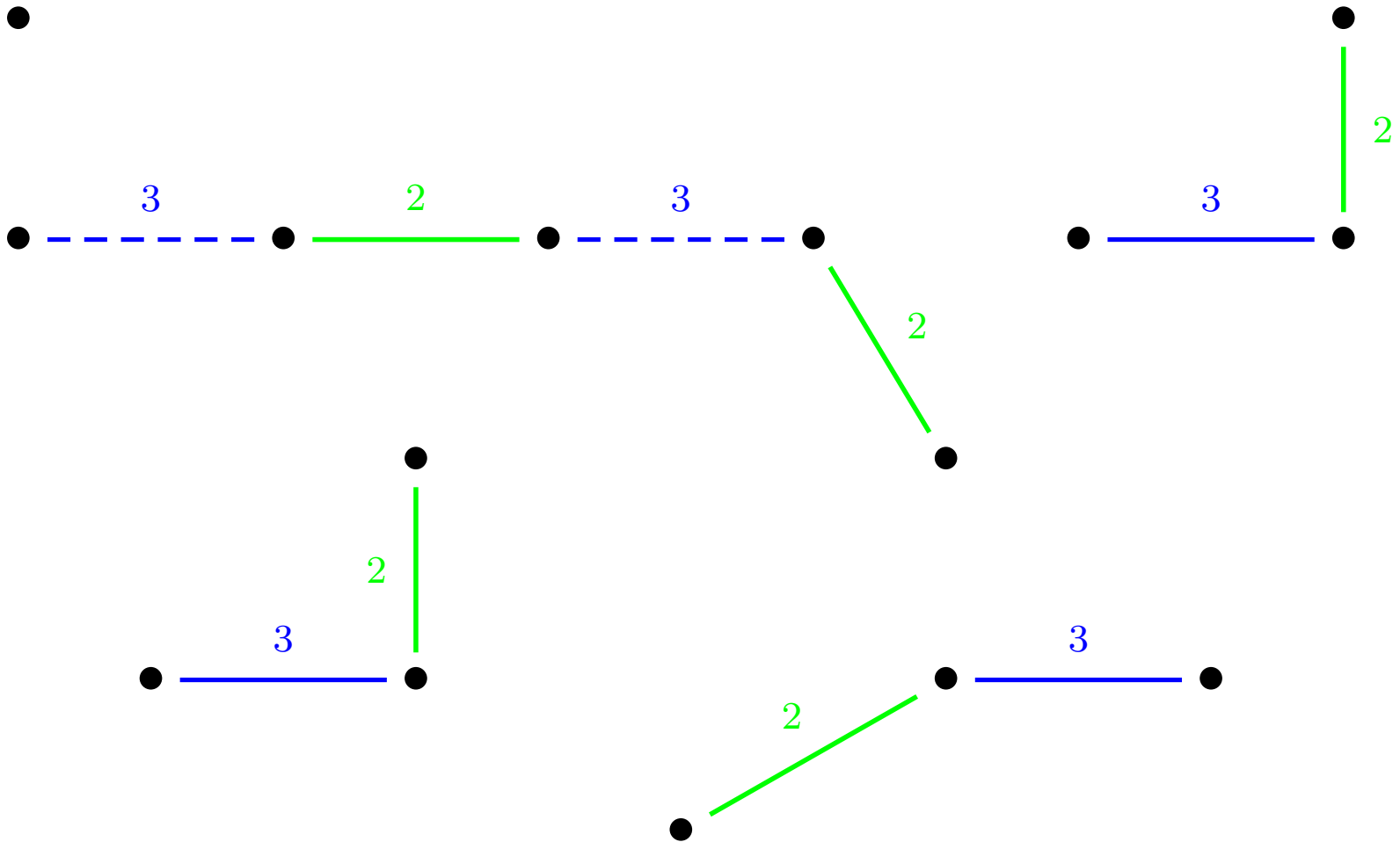
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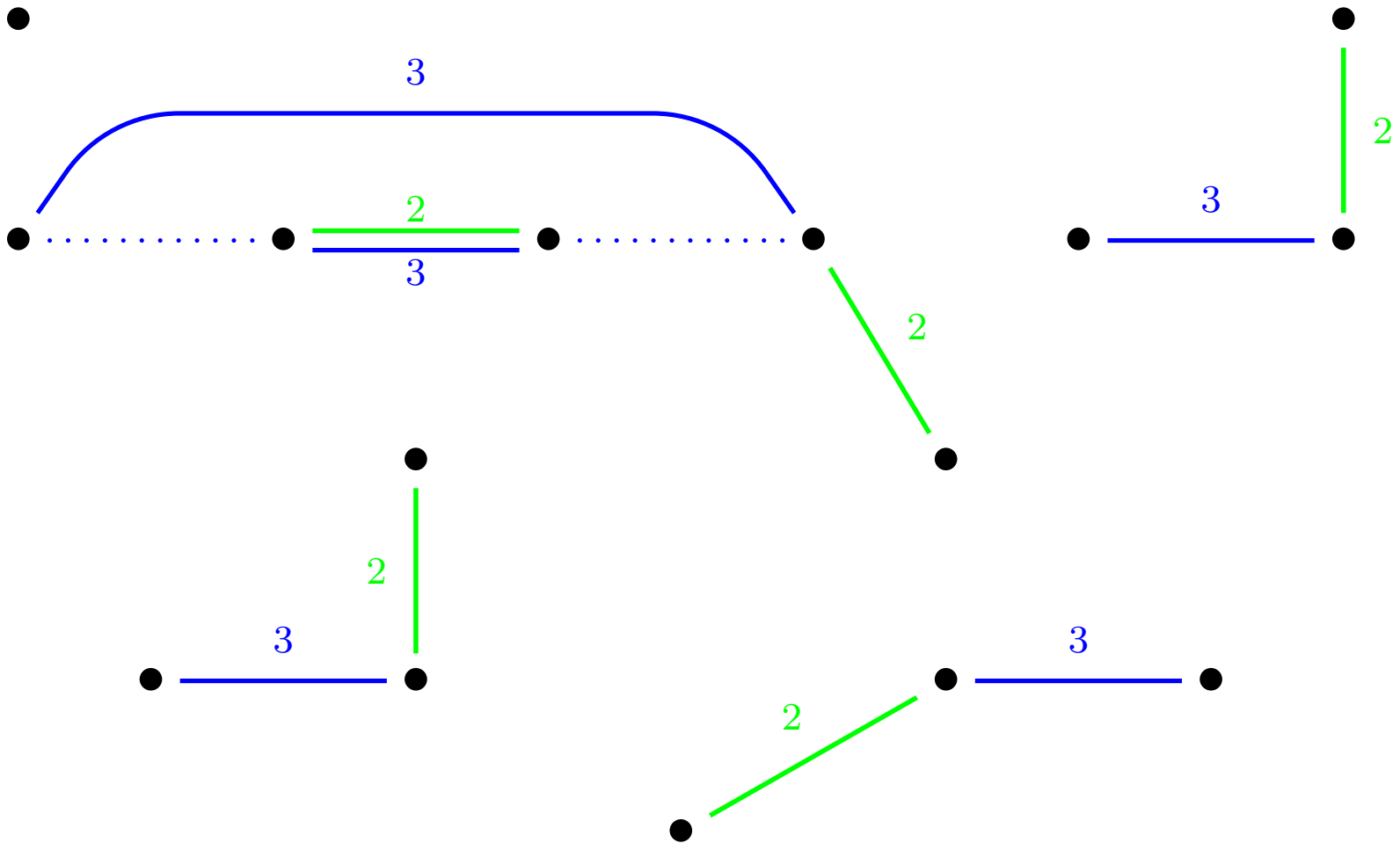
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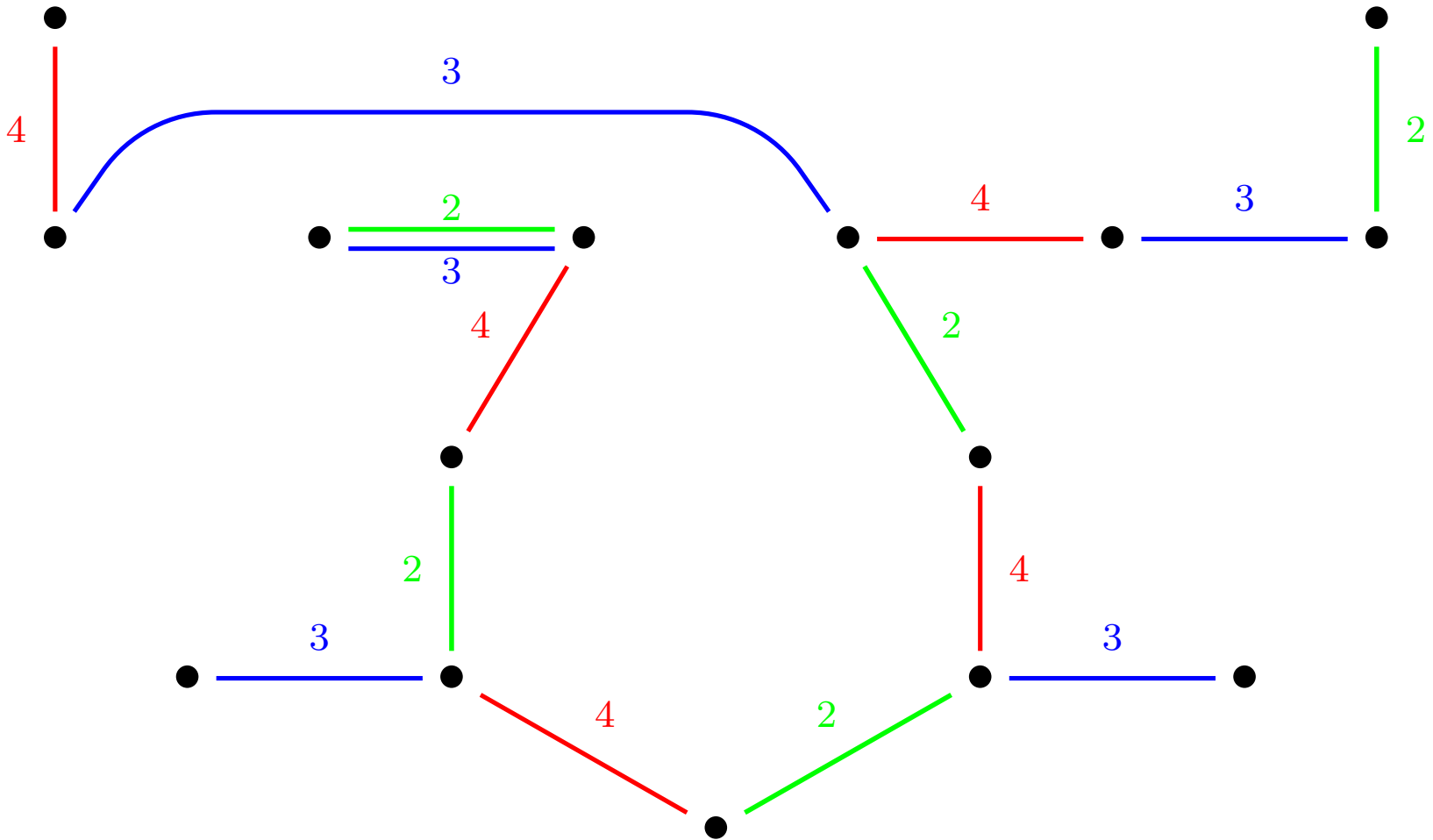
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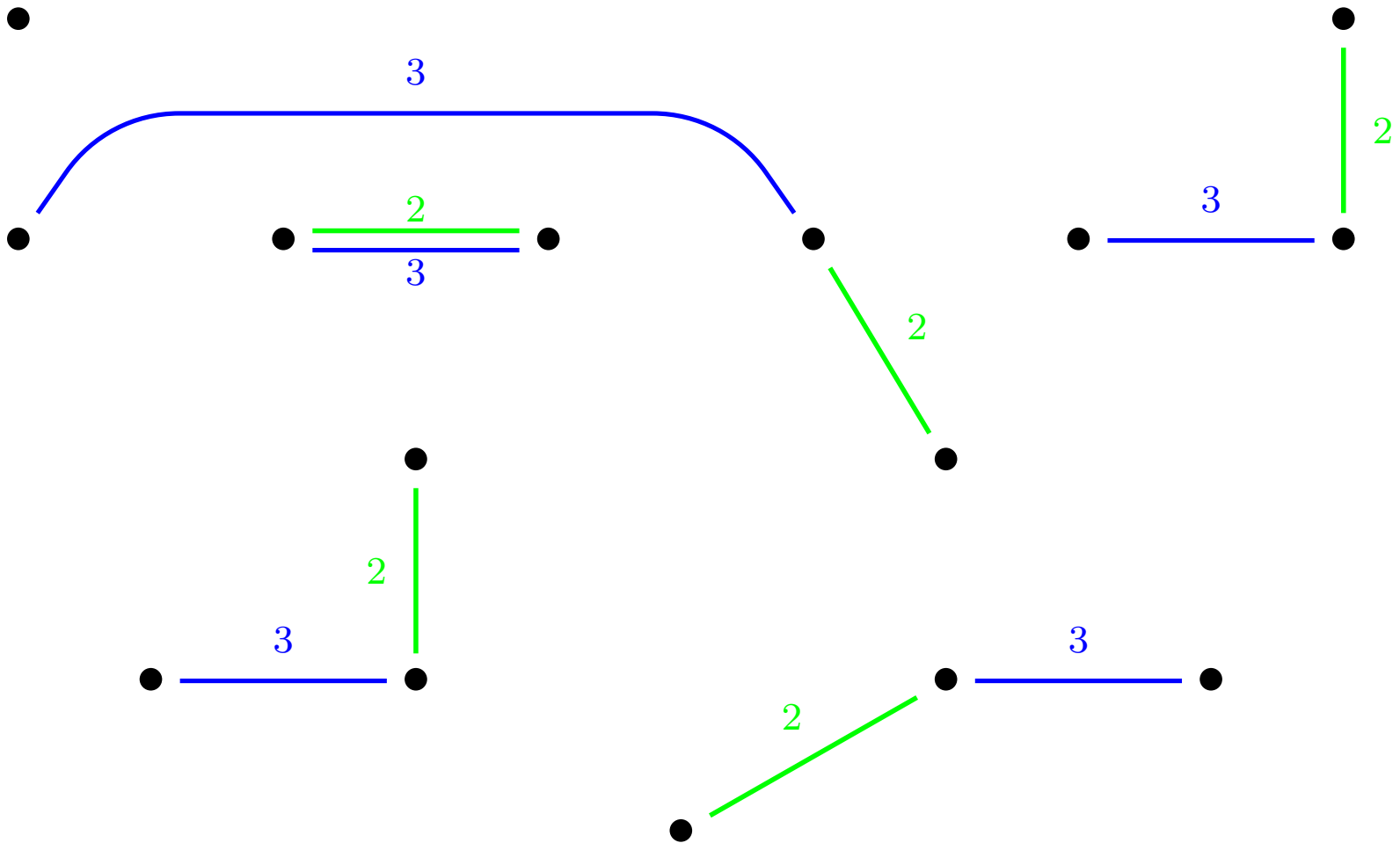
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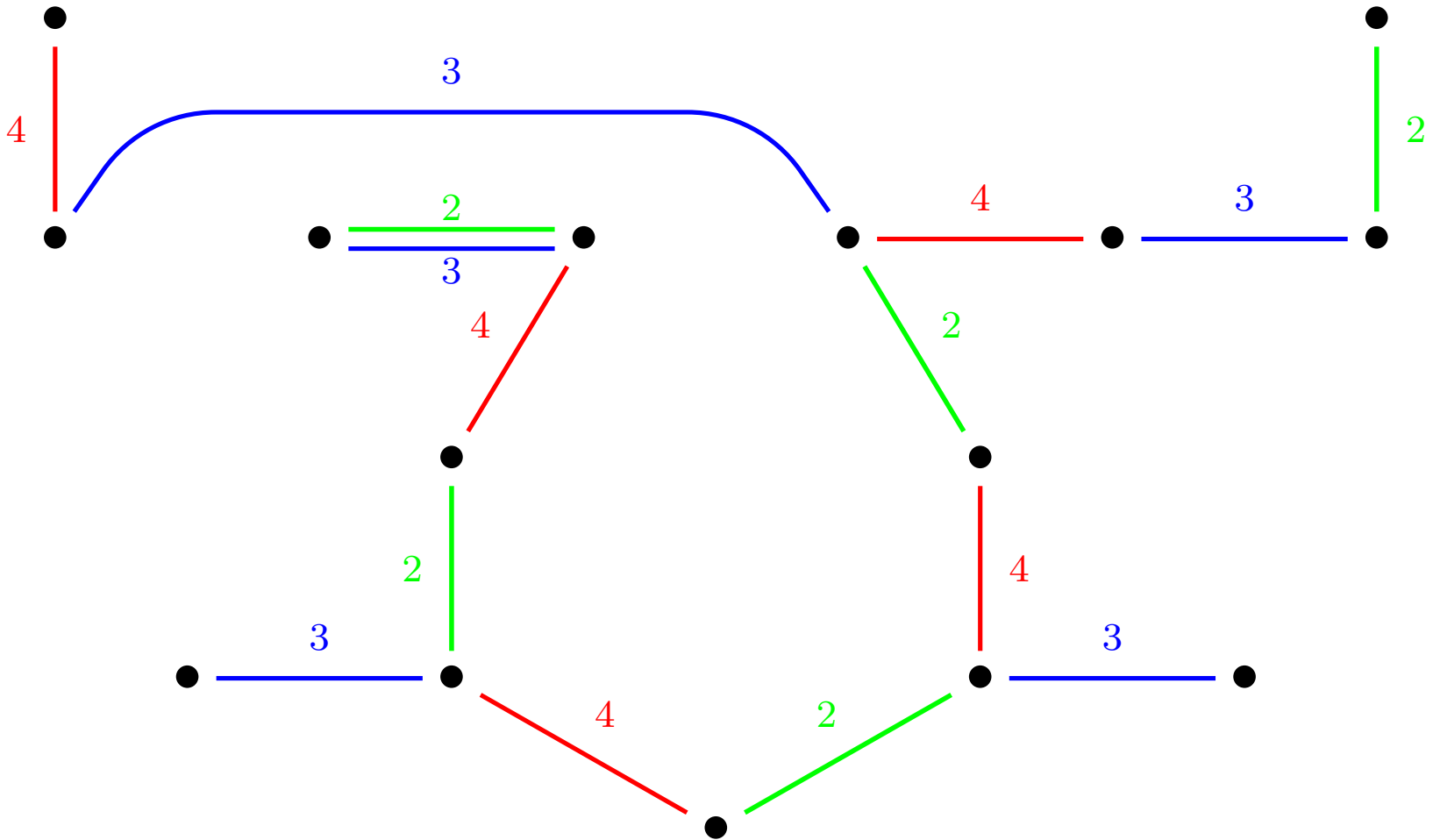
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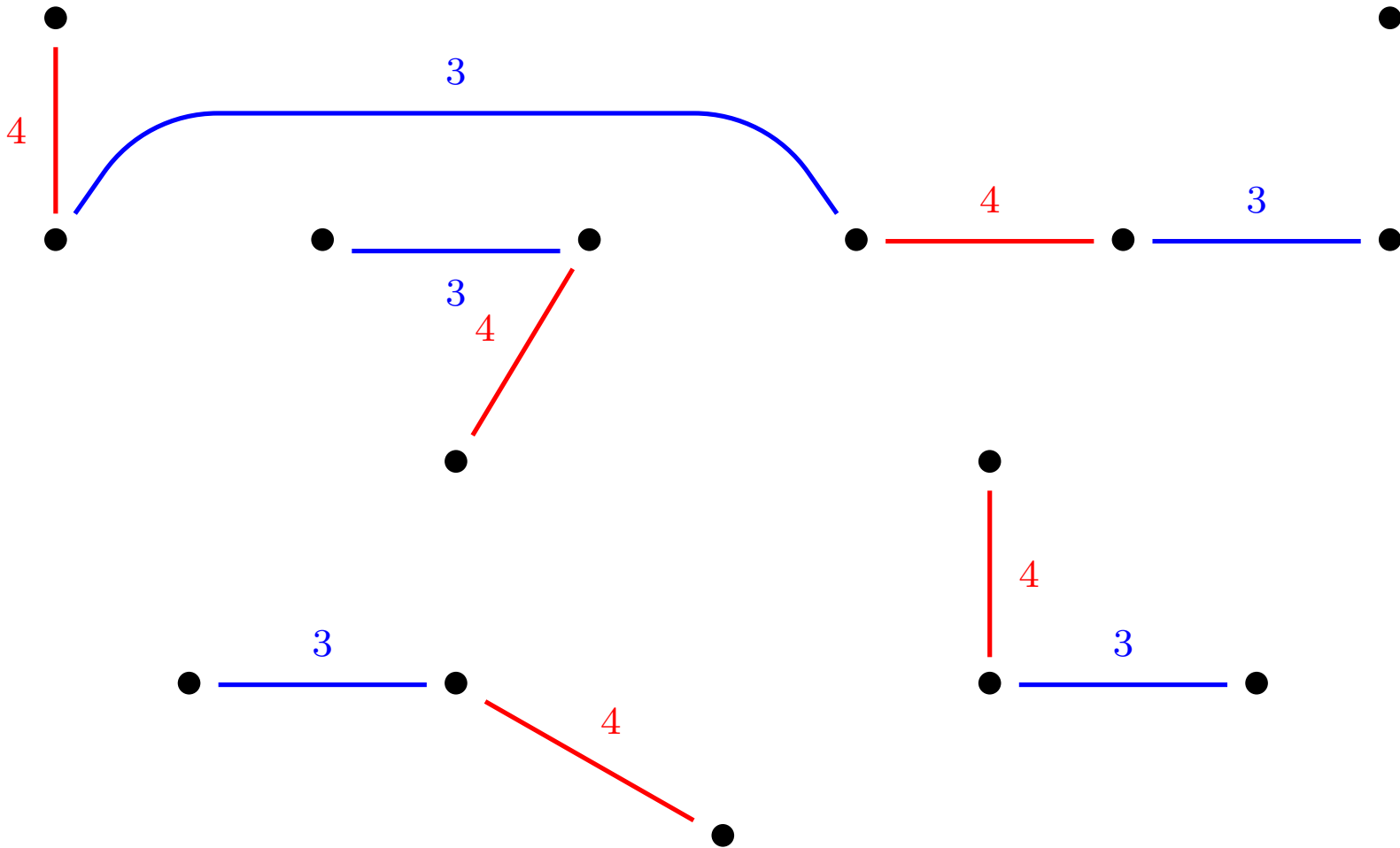
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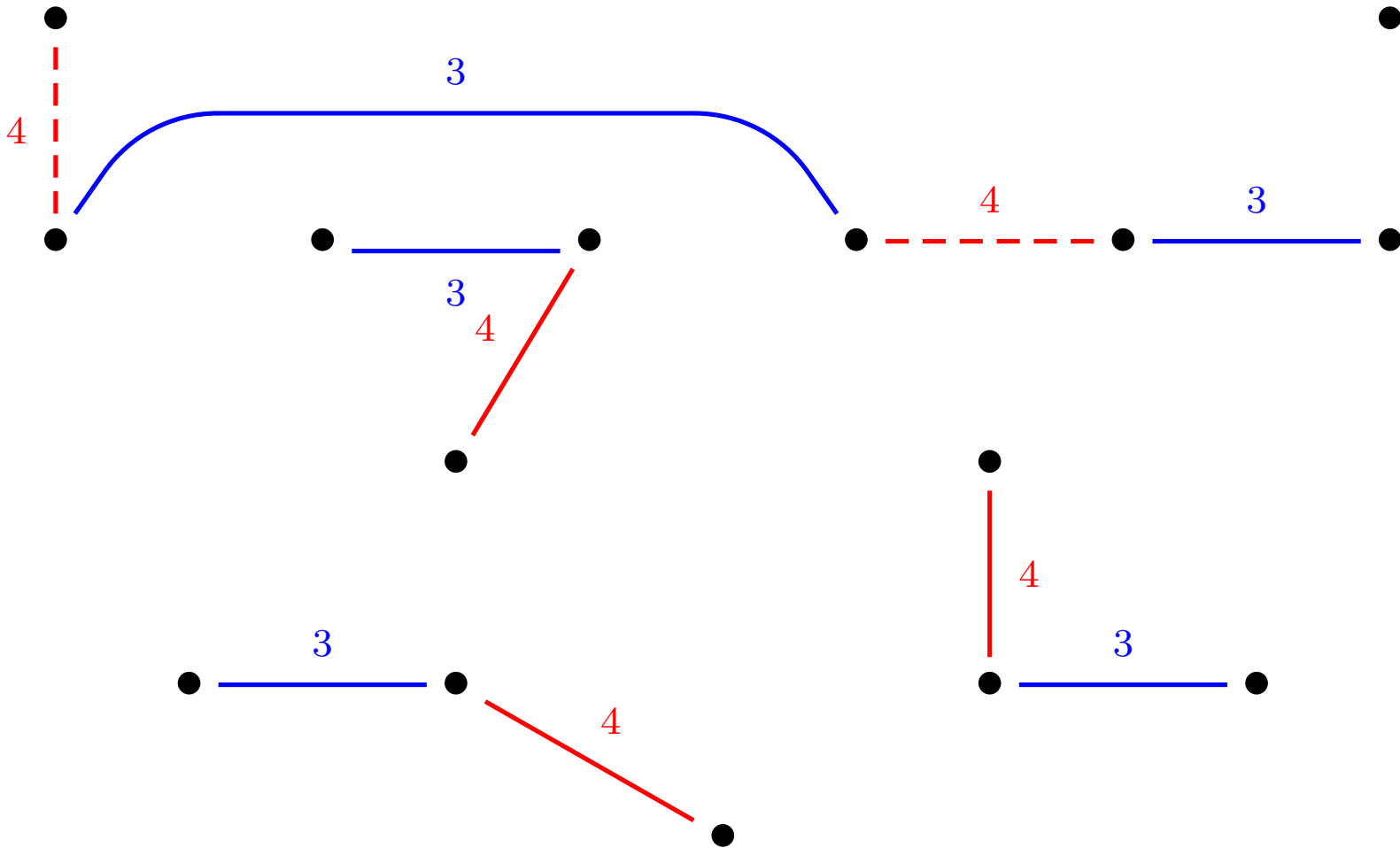
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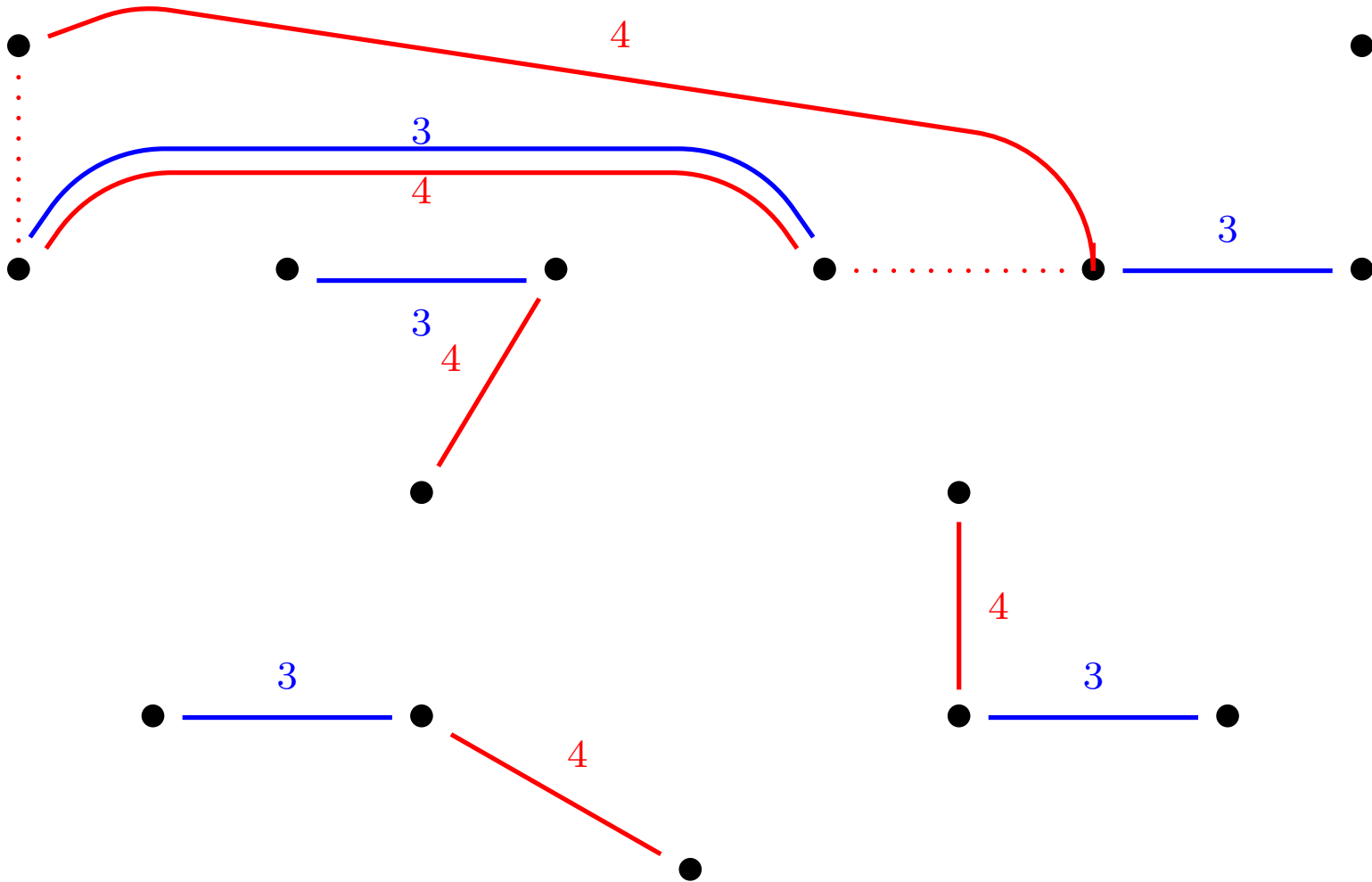
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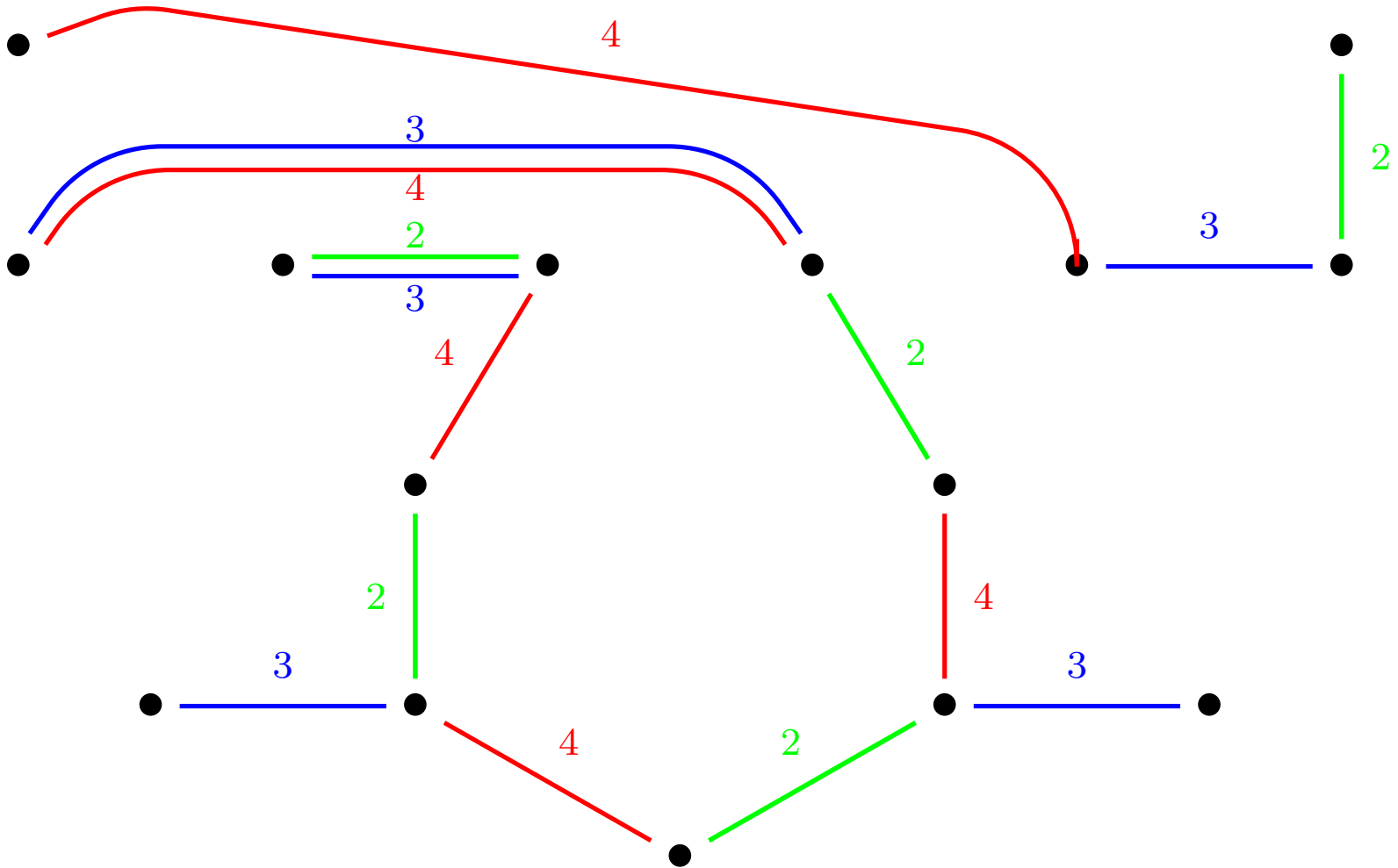
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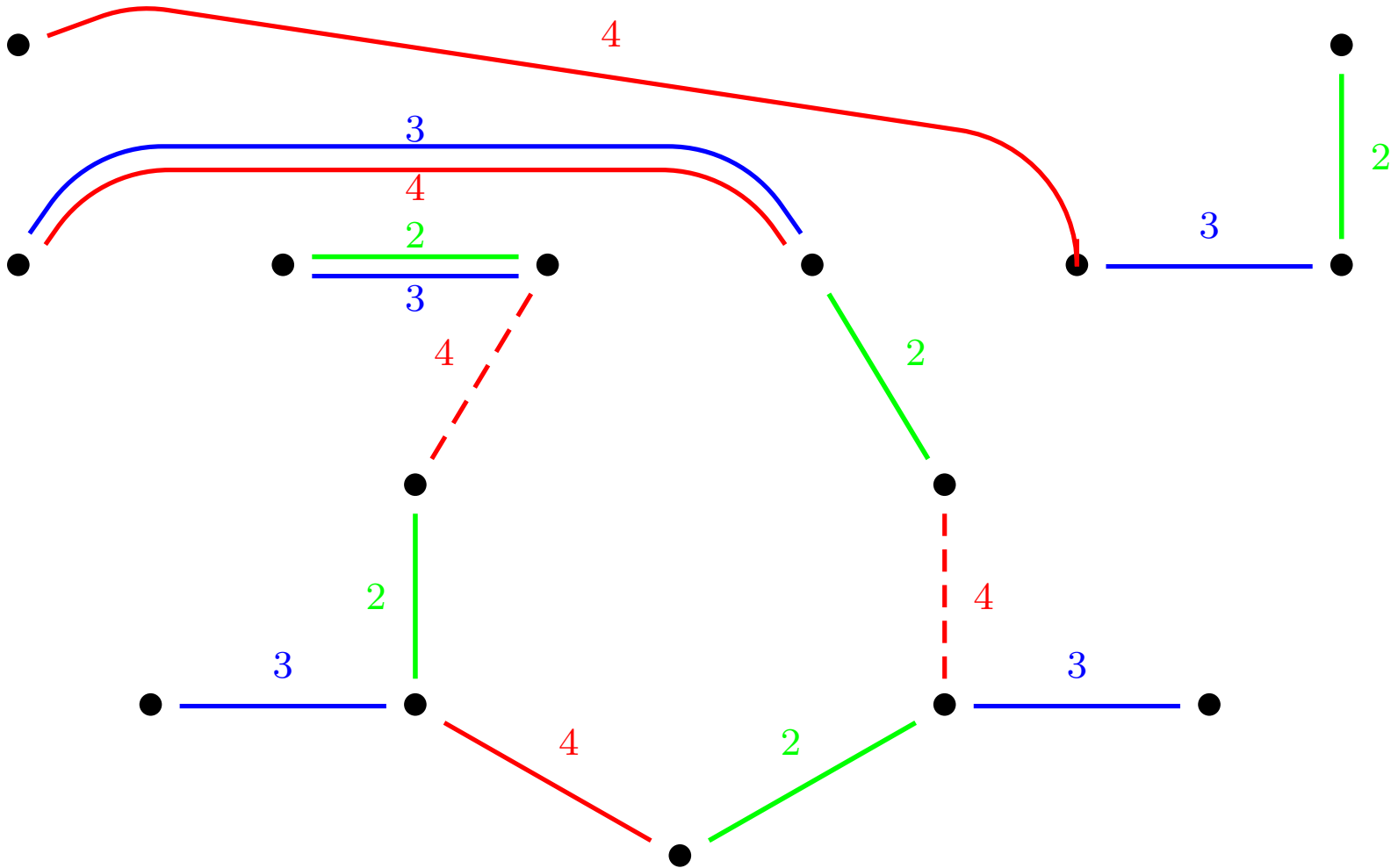
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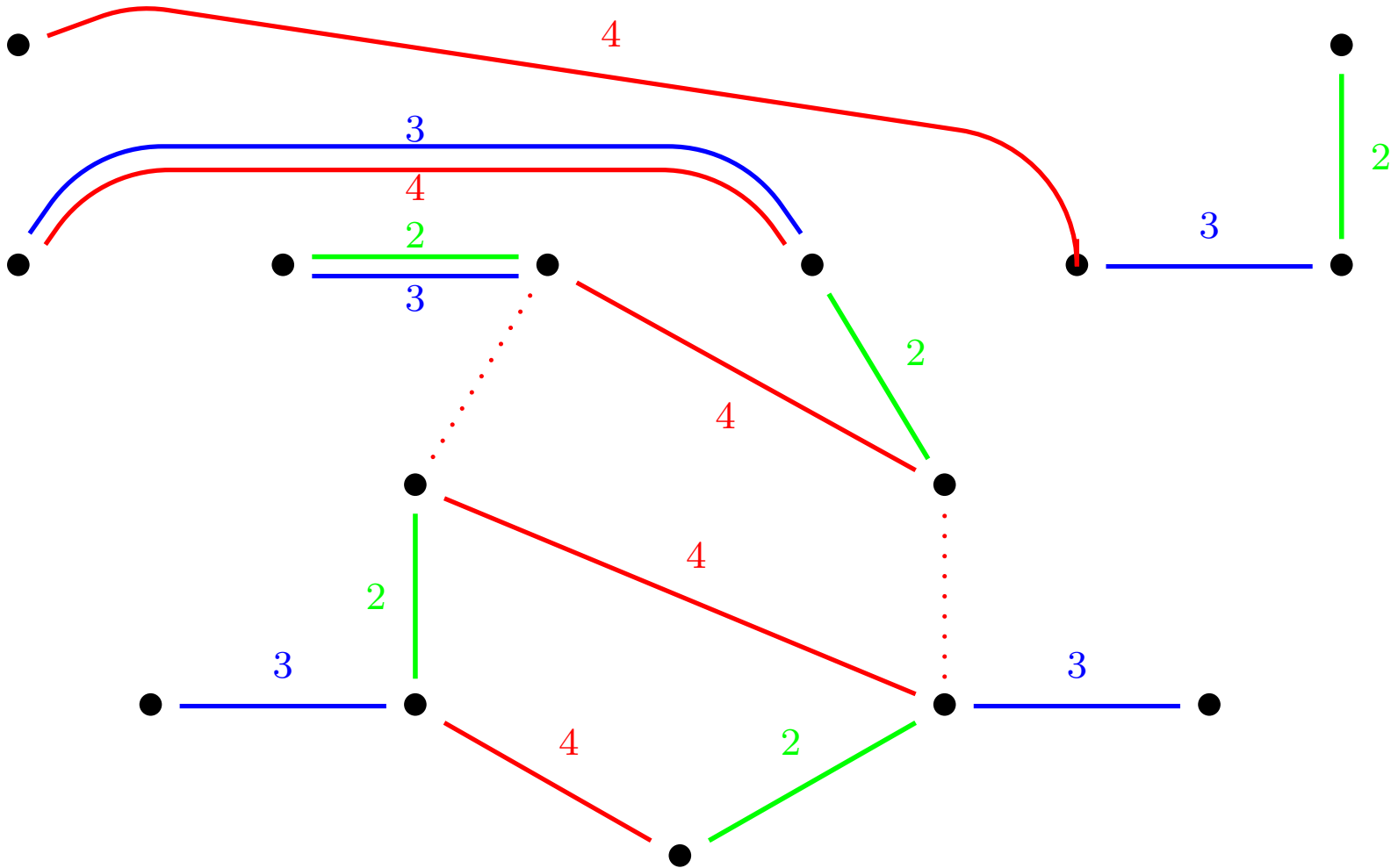
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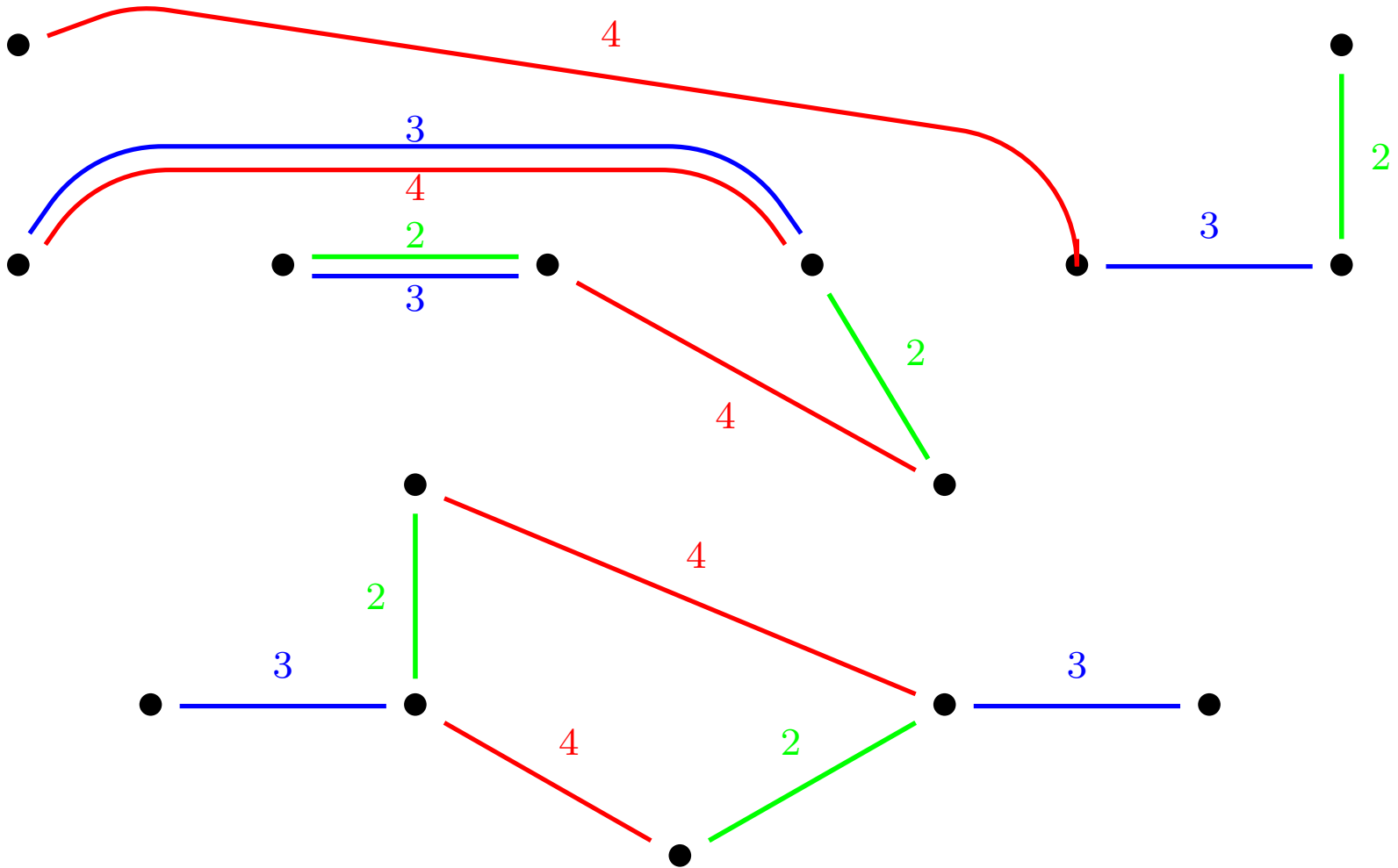
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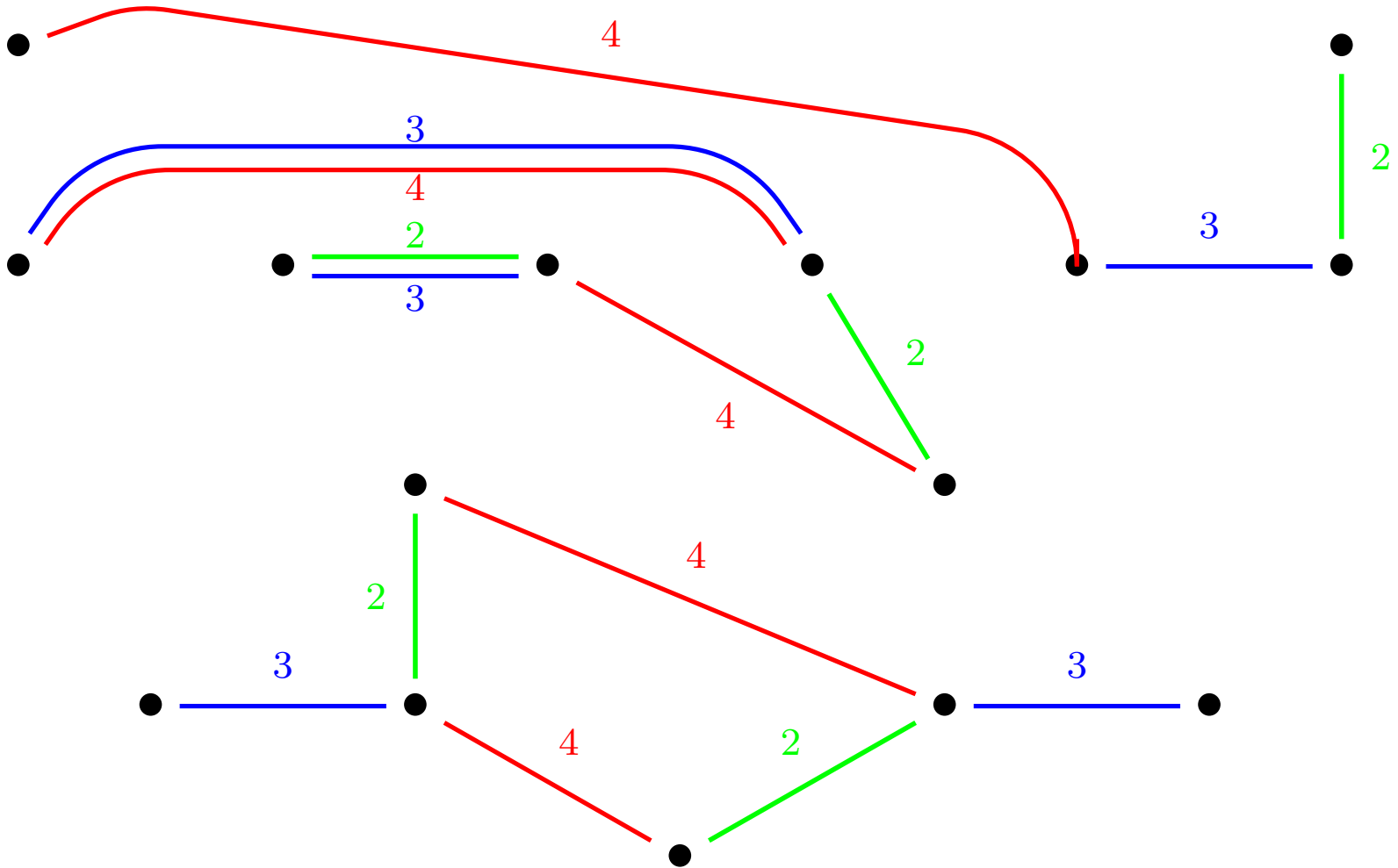
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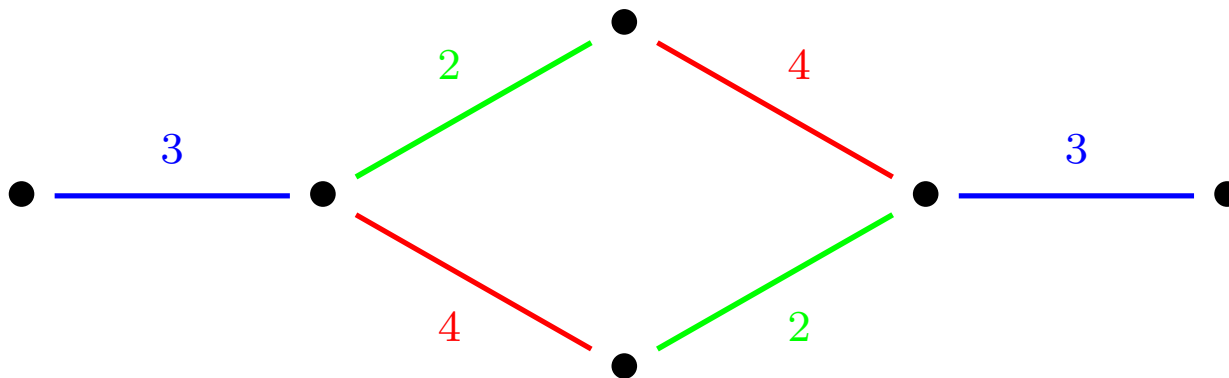
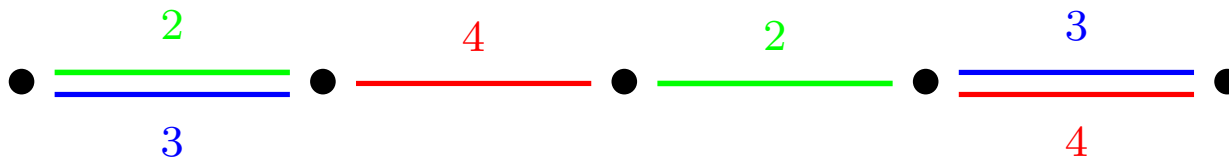
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Macdonald positivity

Corollary. Let $\tilde{\mathcal{H}}_\mu$ be the canonical **DEG** obtained from the **D graph** \mathcal{H}_μ . Then

$$\tilde{K}_{\lambda,\mu}(q, t) = \sum_{\mathcal{C} \cong \mathcal{G}_\lambda} q^{\text{inv}(\mathcal{C})} t^{\text{maj}(\mathcal{C})},$$

where the sum is over connected components of $\tilde{\mathcal{H}}_\mu$ which are isomorphic to the standard **DEG** \mathcal{G}_λ .

In particular, $\tilde{H}_\mu(x; q, t)$ is symmetric and Schur positive.

Preprint available online at

<http://www.math.upenn.edu/~sassaf/PDFs/positivity.pdf>