# Powered by Volcanoes: Three New Algorithms 

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## A 3-volcano of height 2



## $\ell$-volcanoes

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2. For $k>0$, each $v \in V_{k}$ has exactly one neighbor in $V_{k-1}$. All edges not on the surface arise in this manner.
3. For $k<h$, each $v \in V_{k}$ has degree $\ell+1$.

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3. For $k<h$, each $v \in V_{k}$ has degree $\ell+1$.

The integers $\ell, h$, and $\left|V_{0}\right|$ uniquely determine the shape.

## $\ell$-isogenies

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The degree of a (separable) isogeny is $|\operatorname{ker} \phi|$.
We are interested in isogenies of prime degree $\ell$.
Such an isogeny is necessarily cyclic.

The dual isogeny $\hat{\phi}: E_{2} \rightarrow E_{1}$ has the same degree.

## The classical modular polynomial $\Phi_{\ell}$

The polynomial $\Phi_{\ell} \in \mathbb{Z}[X, Y]$ has the property

$$
\Phi_{\ell}\left(j\left(E_{1}\right), j\left(E_{2}\right)\right)=0 \quad \Longleftrightarrow \quad E_{1} \text { and } E_{2} \text { are } \ell \text {-isogenous. }
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Note that $\Phi_{\ell}$ is symmetric in $X$ and $Y$.
The $\ell$-isogeny graph $G_{\ell} / \mathbb{F}_{q}$ has vertex set $\left\{j(E): E / \mathbb{F}_{q}\right\}$ and edges $\left(j_{1}, j_{2}\right)$ whenever $\Phi_{\ell}\left(j_{1}, j_{2}\right)=0\left(\right.$ in $\left.\mathbb{F}_{q}\right)$.

The neighbors of $j$ in $G_{\ell}$ are the roots of $\Phi_{\ell}(X, j) \in \mathbb{F}_{q}[X]$.

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The neighbors of $j$ in $G_{\ell}$ are the roots of $\Phi_{\ell}(X, j) \in \mathbb{F}_{q}[X]$.
$\Phi_{\ell}$ is big: $O\left(\ell^{3} \log \ell\right)$ bits.

## The shape of $G_{\ell}$

An elliptic curve is ordinary (not supersingular) iff its trace is nonzero in $\mathbb{F}_{q}$. Two curves whose $j$-invariants lie in the same component of $G_{\ell}$ are either both ordinary or both supersingular.

## Theorem

The ordinary connected components of $G_{\ell}$ are $\ell$-volcanoes.
(assuming $j \neq 0$, 1728)

Isogenous curves may lie in distint components of $G_{\ell}$. The components of $G_{\ell}$ are a refinement of isogeny classes.

## Finding the floor



## Finding the floor



## Finding the floor



## Finding the floor



## Finding the floor



Finding a shortest path to the floor


Finding a shortest path to the floor


## Finding a shortest path to the floor



## The endomorphism ring End( $E$ )

An endomorphism is an isogeny $\phi: E \rightarrow E$. The multiplication by $m$ map $P \rightsquigarrow m P$ is an example.

The set $\operatorname{End}(E)$ of all endomorphisms of $E$ forms a ring which contains a subring isomorphic to $\mathbb{Z}$.

Over $\mathbb{F}_{q}$ we have $\mathbb{Z} \subsetneq \operatorname{End}(E)$, since

$$
\pi:(X, Y) \rightsquigarrow\left(X^{q}, Y^{q}\right)
$$

is not a multiplication by map.

## End(E) for an ordinary elliptic curve

If $E$ is ordinary than $\operatorname{End}(E) \cong \mathcal{O}$, where $\mathcal{O}$ is an order in an imaginary quadratic field $K$.

We may regard $\pi$ as an element of $\mathcal{O}$ with trace $t$ and norm $q$. The norm equation for $\pi$ has the form

$$
4 q=t^{2}-v^{2} D_{K}
$$

where $K=\mathbb{Q}\left[\sqrt{D_{K}}\right]$ and $v$ is the conductor of $\mathbb{Z}[\pi]$.
We have $\mathbb{Z}[\pi] \subseteq \mathcal{O} \subseteq \mathcal{O}_{K}$, and therefore $\mathcal{O}$ has discriminant $D=u^{2} D_{K}$ for some conductor $u \mid v$.

## The vertical structure of an $\ell$-volcano

Theorem (Kohel)
Let $V_{0}, \ldots, V_{h}$ be the levels of an $\ell$-volcano corresponding to an ordinary component of $G_{\ell} / \mathbb{F}_{q}$.

1. The curves in $V_{i}$ all have the same endomorphism ring type, with discriminant $D_{i}$.
2. $D_{0}$ has conductor prime to $\ell$, and $D_{i}=\ell^{2 i} D_{0}$.

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2. $D_{0}$ has conductor prime to $\ell$, and $D_{i}=\ell^{2 i} D_{0}$.

This implies $\ell^{h} \| v$, allowing us to determine the height.
The endomorphism ring type of an ordinary elliptic curve $E$ is determined by its level on its $\ell$-volcano for each prime $\ell \mid v$.

## The class group action [CM theory]

Suppose $\operatorname{End}(E) \cong \mathcal{O}$, and let $\mathfrak{a}$ an invertible $\mathcal{O}$-ideal. Let $E[a]$ be the points annihilated by all $a \in \mathfrak{a} \subset \mathcal{O} \cong \operatorname{End}(E)$.
There is a separable isogeny $\phi_{\mathfrak{a}}: E \rightarrow E / E[\mathfrak{a}]$ with kernel $E[\mathfrak{a}]$, degree $N(\mathfrak{a})$, and $\operatorname{End}\left(\phi_{\mathfrak{a}}(E)\right) \cong \mathcal{O}$.

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This defines a group action by the ideal group of $\mathcal{O}$ on the set

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\mathcal{E}(\mathcal{O})=\{j(E): \operatorname{End}(E) \cong \mathcal{O}\},
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which factors through the class group $\mathrm{cl}(\mathcal{O})$.

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which factors through the class group $\mathrm{Cl}(\mathcal{O})$.
The above applies over $\mathbb{C}$, but if $E / \mathbb{F}_{q}$ has $\operatorname{End}(E) \cong \mathcal{O}$, then $q$ is the norm of an element of $\mathcal{O}$ and we may reduce to $\mathbb{F}_{q}$.

## The horizontal structure of an ordinary $\ell$-volcano

The degree $d$ of the subgraph on $V_{0}$ is $1+\left(\frac{D_{K}}{\ell}\right)$.
For $d=0$ we have $\left|V_{0}\right|=1$ and for $d=1$ we have $\left|V_{0}\right|=2$.
When $d=2$ there are two $\mathcal{O}$-ideals of norm $\ell, \mathfrak{a}$ and $\overline{\mathfrak{a}}$, and their ideal classes have order $\left|V_{0}\right|$.

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The set $\mathcal{E}(\mathcal{O})$ has size $h(\mathcal{O})$ and is comprised of the surfaces of isomorphic $\ell$-volcanoes corresponding to cosets in $\mathrm{cl}(\mathcal{O})$.

And in general, $\mathcal{E}(\mathcal{O})$ is a torsor for $\mathrm{cl}(\mathcal{O})$.

## The CM method

If $E / \mathbb{F}_{q}$ has $N=q+1-t$ points, with $t \not \equiv 0$ in $\mathbb{F}_{q}$, then

$$
4 q=t^{2}-v^{2} D
$$

where $D$ is the discriminant of $\mathcal{O} \cong \operatorname{End}(E)$. Conversely, any curve with $\operatorname{End}(E) \cong O$ has trace $\pm t$.

The Hilbert class polynomial $H_{D} \in \mathbb{Z}[X]$ is defined by

$$
H_{D}(X)=\prod_{j \in \mathcal{E}(\mathcal{O})}(X-j)
$$

Its roots are the $j$-invariants of curves with $\operatorname{End}(E) \cong O$.
Given a root of $H_{D}$ in $\mathbb{F}_{q}$, we may construct $E / \mathbb{F}_{q}$ with $N$ points.

## Computing $H_{D}(X)$ with the CRT [ALV '06, BBEL '08]

To compute $H_{D} \in \mathbb{F}_{q}[X]$ it suffices to compute $H_{D}$ modulo many "small" primes $p$ and apply the Chinese Remainder Theorem.

For primes of the form $4 p=t_{p}^{2}-v_{p}^{2} D, H_{D}$ splits completely over $\mathbb{F}_{p}$ and we may compute $H_{D} \bmod p$ by finding its roots.

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To find the first root, generate random curves over $\mathbb{F}_{p}$ until we find one with $\operatorname{End}(E) \cong \mathcal{O}$ (or any $E$ with trace $\pm t$ ).

To enumerate the other roots, use the group action of $\mathrm{cl}(\mathcal{O})$.

## Improvements [S '09]

The CRT approach to computing $H_{D}$ can be improved:

1. Compute $H_{D} \bmod P$ in $O\left(|D|^{1 / 2+\epsilon} \log P\right)$ space.
2. Generate "random" curves with prescribed torsion.
3. Make $v_{p}$ large (bigger volcanoes are easier to find).
4. Use an optimal presentation of $\mathrm{cl}(\mathcal{O})$ to minimize norms.

## An example of a polycyclic presentation

For $D=-79947, \mathrm{cl}(D)$ is cyclic of order $h(D)=100$.
It is generated by the class of an ideal with norm 19.
But $\mathrm{cl}(\boldsymbol{D})$ is also generated by classes $\alpha_{2}$ and $\alpha_{13}$ of ideals of norm 2 and 13. The elements $\alpha_{2}$ and $\alpha_{13}$ have orders 20 and 50 and are not independent $\left(\alpha_{13}^{5}=\alpha_{2}^{18}\right)$.

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Nevertheless, every $\beta \in \operatorname{cl}(D)$ can be written uniquely as

$$
\beta=\alpha_{2}^{e_{2}} \alpha_{13}^{e_{13}}
$$

with $0 \leq e_{2}<20$ and $0 \leq e_{13}<5$.
Using this presentation is about 100 times faster.

## Running the rim



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## Record-breaking CM constructions

Largest |D|
Old Record (June 2008, complex analytic [Enge])

$$
D=-70,901,505,867 \quad h(D)=51,244
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New Record (October 2008, CRT method [Enge-S])

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D=-102,197,306,669,747 \quad h(D)=2,014,236
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Largest $h(D)$
Old Record (January 2006, complex analytic [Enge])

$$
D=-2,093,236,031 \quad h(D)=100,000
$$

New Record (April 2009, CRT method, [Bröker-S])

$$
D=-4,058,817,012,071 \quad h(D)=5,000,000
$$

## Performance comparison

| -D | $h(D)$ | Analytic $\mathfrak{w}_{3,13}$ |  | CRT $\mathfrak{f}^{2}$ |  | CRT $f$ |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | height | time | height* | time* | height | time |
| 6961631 | 5000 | 9.5 k | 28 | 9.5 k | 4.9 | 3.8 k | 2.0 |
| 23512271 | 10000 | 20k | 210 | 20k | 24 | 8.0k | 9.1 |
| 98016239 | 20000 | 45k | 1,800 | 45k | 120 | 18k | 46 |
| 357116231 | 40000 | 97k | 14,000 | 97k | 574 | 38k | 220 |
| 2093236031 | 100000 | 265k | 260,000 | 265k | 4,400 | 103k | 1,600 |

## Complex Analytic vs. CRT method

(2.4 GHz AMD Opteron CPU seconds)
*increased to match the height bound for $\mathfrak{w}_{3,13}$.

## Computing End(E) [Bisson-S '09]

Given $E / \mathbb{F}_{q}$ we may compute $t$ and factor $4 q-t^{2}$ to obtain

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The discriminant of $\operatorname{End}(E) \cong \mathcal{O}$ is $D=u^{2} D_{K}$ for some $u \mid v$ To determine $\operatorname{End}(E)$ it suffices to compute $u$.

Let $u_{1}, \ldots, u_{n}$ be the factors of $v$. To distinguish $u$, we seek relations that hold in some $\operatorname{cl}\left(u_{i}^{2} D_{K}\right)$ but not others.

We test these relations in the isogeny graph by walking along the surface of various $\ell$-volcanoes.

## Relations in class groups

A relation $R$ is a pair of vectors $\left(\ell_{1}, \ldots, \ell_{r}\right)$ and $\left(e_{1}, \ldots, e_{r}\right)$, with $\ell_{i} \nmid v$ and $\left(\frac{D_{K}}{\ell_{i}}\right)=1$.

We say $R$ holds in $\mathrm{cl}(D)$ if for each $i$ there is an $\alpha_{i} \in \mathrm{cl}(D)$ containing an ideal of norm $\ell_{i}$ such that $\alpha_{1}^{e_{1}} \cdots \alpha_{r}^{e_{r}}=1$.

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We say $R$ holds in $\mathrm{cl}(D)$ if for each $i$ there is an $\alpha_{i} \in \mathrm{cl}(D)$ containing an ideal of norm $\ell_{i}$ such that $\alpha_{1}^{e_{1}} \cdots \alpha_{r}^{e_{r}}=1$.
More generally, define the cardinality of $R$ in $\mathrm{cl}(D)$ by

$$
\# R / \operatorname{cl}(D)=\#\left\{\tau \in\{ \pm 1\}^{r}: \prod \alpha_{i}^{\tau_{i} e_{i}}=1 \operatorname{incl}(D)\right\}
$$

For $p \mid v$, let $D_{1}=(v / p)^{2} D_{K}$ and $D_{2}=p^{2} D_{K}$. We want

$$
\# R / \operatorname{cl}\left(D_{1}\right)>\# R / \operatorname{cl}\left(D_{2}\right)
$$

## Counting relations in the isogeny graph

To compute $\# R / \mathrm{cl}(\mathcal{O})$ :

1. Let $J_{0}$ be a list consisting of $j(E)$.
2. For $i$ from 1 to $r$ :

- For each $j \in J_{i-1}$, walk $e_{i}$ steps in both directions on the surface of the $\ell_{i}$-volcano and append the endpoints to $J_{i}$.

3. Output the number of times $j(E)$ occurs in $J_{r}$.

## Counting relations in the isogeny graph

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3. Output the number of times $j(E)$ occurs in $J_{r}$.

To compute $\# R / \mathrm{cl}(\mathcal{O})$ efficiently, we use smooth relations, where $\ell_{i}, e_{i}$, and $r$ are all small.

## Record-breaking End( $E$ ) computations

Heuristically, we achieve a running time of $L[1 / 2, \sqrt{3} / 2]$.
Over a 200-bit prime field, under 15 minutes.
Over a 256 -bit prime field, about 4 hours.
These are worst-case examples (average case is easy).

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These are worst-case examples (average case is easy).
Kohel's algorithm has complexity $O\left(q^{1 / 3}\right)$ (under the GRH). It cannot feasibly compute End( $E$ ) over a cryptographic size field when $v$ contains a large prime factor.

## Computing $\Phi_{\ell}$ with the CRT method [Bröker-Lauter-S]

Choose CRT primes $p \equiv 1 \bmod \ell$ with $4 p=t^{2}-v^{2} \ell^{2} D$. Suppose we have an $\ell$-volcano of height 1 with $\left|V_{0}\right| \geq \ell+2$. (we may pick $D$ to ensure this).

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We can "construct" this volcano without using $\Phi_{\ell}$ :

1. Use $H_{D}(X)$ to find the surface.
2. Apply the action of $\mathrm{cl}(D)$ to enumerate the surface.
3. Use Velu's formula to descend to the floor.
4. Apply the action of $\mathrm{cl}\left(\ell^{2} D\right)$ to enumerate the floor.

From this we can interpolate $\Phi_{\ell} \bmod p$.

## Record-breaking $\Phi_{\ell}$ computations

The time to compute $\Phi_{\ell}$ is $O\left(\ell^{3} \log ^{3+\epsilon} \ell\right)$ [GRH]. Faster than the best alternative by a factor of $\log \ell$.

Record $\Phi_{\ell}$ computations (classical)
Computed $\Phi_{\ell}$ for all $\ell<3000$, and up to $\ell=5003$.
Output is generated at a rate of about $5 \mathrm{Mb} / \mathrm{s}$.
Previous record: $\ell<360$ [Rubinstein-Seroussi].

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Record modular polynomial computations (Weber f)
Computed $\Phi_{\ell}$ for all $\ell<10000$ and up to $\ell=50021$.
Preprint in preparation.

## Modular polynomials for $\ell=7$

Classical:

$$
\begin{aligned}
& X^{8}+Y^{8}-X^{7} Y^{7}+5208 X^{7} Y^{6}-10246068 X^{7} Y^{5}+9437674400 X^{7} Y^{4}-4079701128594 X^{7} Y^{3}+ \\
& 720168419610864 X^{7} Y^{2}-34993297342013192 X^{7} Y+104545516658688000 X^{7}+ \\
& \ldots(2 \text { pages omitted }) \ldots+
\end{aligned}
$$

$$
13483958224762213714698012883865296529472356352000000000000000 Y^{3}+
$$

$$
1464765079488386840337633731737402825128271675392000000000000000000 Y^{2}
$$

Atkin:

$$
\begin{aligned}
& X^{8}-X^{7} Y+744 X^{7}+196476 X^{6}+357 X^{5} Y+21226520 X^{5}+1428 X^{4} Y+ \\
& 803037606 X^{4}-31647 X^{3} Y+14547824088 X^{3}-204792 X^{2} Y+138917735740 X^{2}+ \\
& 186955 X Y+677600447400 X+Y^{2}+2128500 Y+1335206318625
\end{aligned}
$$

Canonical:

$$
X^{8}+28 X^{7}+322 X^{6}+1904 X^{5}+5915 X^{4}+8624 X^{3}+4018 X^{2}-X Y+748 X+49
$$

Weber:

$$
x^{8}+y^{8}-x^{7} y^{7}+7 X^{4} y^{4}-8 X Y
$$

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