## Modular polynomials and isogeny volcanoes

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#### Isogenies

An *isogeny*  $\phi : E_1 \rightarrow E_2$  is a morphism of elliptic curves, a nonzero rational map that preserves the identity.

Over a finite field,  $E_1$  and  $E_2$  are isogenous if and only if  $\#E_1(\mathbb{F}_q) = \#E_2(\mathbb{F}_q).$ 

# Some applications of isogenies

Isogenies make hard problems easier:

- Counting the points on E.
  Polynomial time (Schoof-Elkies-Atkin).
- Constructing *E* with the CM method.  $|D| > 10^{15}$  (BBEL, S, Enge-S).
- Computing the endomorphism ring of *E*. Heuristically subexponential time (Bisson-S).

These algorithms all rely on modular polynomials  $\Phi_{\ell}(X, Y)$ .

# Properties of isogenies

#### Degree

The kernel of  $\phi : E_1 \to E_2$  is a finite subgroup of  $E_1(\overline{F})$ . When  $\phi$  is separable, we have  $|\ker \phi| = \deg \phi$ .

An  $\ell$ -isogeny is a (separable) isogeny of degree  $\ell$ . For prime  $\ell$ , the kernel is necessarily cyclic.

Orientation We say that  $\phi : E_1 \to E_2$  is *horizontal* if  $End(E_1) \cong End(E_2)$ . Otherwise  $\phi$  is *vertical*.

## Isogenies from kernels

Any finite subgroup G of  $E(\overline{F})$  determines a separable isogeny with G as its kernel

Given G, we can compute  $\phi$  explicitly using Vélu's formula.

The complexity depends both on the size of ker  $\phi$ , and the field in which the points of ker  $\phi$  are defined.

If *E* is defined over *F*, so is  $\phi$ , but the points in ker  $\phi$  may have coordinates in an extension of degree up to  $\ell^2 - 1$ .

## The classical modular polynomial $\Phi_\ell$

The modular function  $j : \mathbb{H} \to \mathbb{C}$  is a complex analytic function

 $j(z) = 1/q + 744 + 196884q + 21493760q^2 + \dots,$ 

where  $q = e^{2\pi i z}$ . The function  $j(\ell z)$  is algebraic over  $\mathbb{C}(j)$ , and its minimal polynomial  $\Phi_{\ell}(X)$  has coefficients in  $\mathbb{Z}[j]$ .

$$\Phi_{\ell}(X, Y) = \Phi_{\ell}(Y, X);$$
  $\deg_X \Phi_{\ell} = \ell + 1.$ 

The modular equation  $\Phi_{\ell}(X, Y) = 0$  parameterizes pairs of elliptic curves related by a cyclic  $\ell$ -isogeny.

## Parametrizing isogenies

Assuming char  $F \neq \ell$ , for all elliptic curves  $E_1/F$  and  $E_2/F$ :

 $\Phi_{\ell}(j(E_1), j(E_2)) = 0 \iff E_1 \text{ and } E_2 \text{ are } \ell \text{-isogenous.}$ 

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The  $\ell$ -isogeny graph  $G_{\ell}$  has vertex set  $\{j(E) : E/F\}$ , and edges  $(j_1, j_2)$  whenever  $\Phi_{\ell}(j_1, j_2) = 0$ .

The neighbors of  $j_0$  are the roots of  $\Phi_{\ell}(X, j_0)$  that lie in *F*.

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 $\Phi_{\ell}$  is big:  $O(\ell^3 \log \ell)$  bits.

l	coefficients	largest	average	total
127	8258	7.5kb	5.3kb	5.5MB
251	31880	16kb	12kb	48MB
503	127262	36kb	27kb	431MB
1009	510557	78kb	60kb	3.9GB
2003	2009012	166kb	132kb	33GB
3001	4507505	259kb	208kb	117GB
4001	8010005	356kb	287kb	287GB
5003	12522512	454kb	369kb	577GB
10007	50085038	968kb	774kb	4.8TB
20011	200250080	2.0Mb	1.6Mb	40TB*

Size of  $\Phi_{\ell}(X, Y)$ 

\*Estimated

## Algorithms to compute $\Phi_\ell$

#### q-expansions:

 $\begin{array}{ll} (\text{Atkin ?, Elkies '92, '98, LMMS '94, Morain '95, Müller '95, BCRS '99)} \\ \Phi_{\ell}: & O(\ell^4 \log^{3+\epsilon} \ell) & (\text{via the CRT}) \\ \Phi_{\ell} \bmod p: & O(\ell^3 \log \ell \log^{1+\epsilon} p) & (p > \ell+1) \end{array}$ 

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evaluation-interpolation: (Enge 2009) $\Phi_{\ell}$ : $O(\ell^3 \log^{4+\epsilon} \ell)$ (floating-point) $\Phi_{\ell}$  mod m: $O(\ell^3 \log^{4+\epsilon} \ell)$ (reduces  $\Phi_{\ell}$ )

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Under the GRH, we find many such p with log  $p = O(\log \ell)$ .

 $\begin{array}{ll} \Phi_{\ell} : & O(\ell^3 \log^{3+\epsilon} \ell) & (\text{via the CRT}) \\ \Phi_{\ell} \bmod m : & O(\ell^3 \log^{3+\epsilon} \ell) & (\text{via the explicit CRT}) \end{array}$ 

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In practice the algorithm is much faster than other methods. It is probabilistic, but its output is unconditionally correct.

## **Explicit Chinese Remainder Theorem**

Suppose  $c \equiv c_i \mod p_i$  for distinct primes  $p_i, \ldots, p_n$ . Then

 $c\equiv \sum c_i a_i M_i \mod M,$ 

where  $M = \prod p_i$ ,  $M_i = M/p_i$  and  $a_i = 1/M_i \mod p_i$ .

With M > 4c, the explicit CRT computes  $c \mod m$  directly via

$$c \equiv \left(\sum c_i a_i M_i - rM\right) \mod m,$$

where the integer  $r \approx \sum a_i c_i / p_i$  (use  $O(\log n)$  bits of precision).

Using an online algorithm, this can be applied to *N* coefficients *c* in parallel, using  $O(\log M + n \log m + N(\log m + \log n))$  space.

Montgomery-Silverman, Bernstein, S.

# Some performance highlights

#### Level records

- **1. 5003**: Φ<sub>ℓ</sub>
- **2. 20011**: Φ<sub>ℓ</sub> mod *m*
- **3.** 60013: Φ<sup>f</sup><sub>ℓ</sub>

#### Speed records

1. 251:  $\Phi_{\ell}$  in 28s $\Phi_{\ell}$  mod m in 4.8s(vs 688s)2. 1009:  $\Phi_{\ell}$  in 2830s $\Phi_{\ell}$  mod m in 265s(vs 107000s)

Single core CPU times (AMD 3.0 GHz), using  $m \approx 2^{256}$ .

Effective throughput when computing  $\Phi_{1009} \mod m$  is over 100 Mb/s.



## A 3-volcano of depth 2



#### *ℓ*-volcanoes

An  $\ell$ -volcano is a connected undirected graph whose vertices are partitioned into levels  $V_0, \ldots, V_d$ , such that:

- 1. The subgraph on  $V_0$  (the *surface*) is a regular connected graph of degree at most 2.
- 2. For i > 0, each  $v \in V_i$  has exactly one neighbor in  $V_{i-1}$ . All edges not on the surface arise in this manner.
- 3. For i < d, each  $v \in V_i$  has degree  $\ell$ +1.

The integers  $\ell$ , d, and  $|V_0|$  uniquely determine the shape.

# The $\ell$ -isogeny graph $G_{\ell}$

Some facts about  $G_{\ell}$  (Kohel, Fouquet-Morain):

- ► The ordinary components of G<sub>ℓ</sub> are ℓ-volcanoes (provided they don't contain j = 0, 1728).
- ► The curves in level V<sub>i</sub> of a given ℓ-volcano all have the same endomorphism ring, isomorphic to an imaginary quadratic order O<sub>i</sub>.
- The order  $\mathcal{O}_0$  is maximal at  $\ell$ , and  $[\mathcal{O}_0 : \mathcal{O}_i] = \ell^i$ .

Curves in the same  $\ell$ -volcano are necessarily isogenous, but isogenous curves need not lie in the same  $\ell$ -volcano.

Let  $E/\mathbb{F}_q$  be an ordinary elliptic curve with  $End(E) \cong \mathcal{O}$ . The class group  $cl(\mathcal{O})$  acts on the set

 $\{j(E/\mathbb{F}_q): \operatorname{End}(E) \cong \mathcal{O}\}.$ 

Horizontal  $\ell$ -isogenies are the action of an ideal with norm  $\ell$ .

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A horizontal isogeny of large degree may be equivalent to a sequence of isogenies of small degree, via relations in cl(O).

Under the ERH this is always true, and "small" =  $O(\log^2 |D|)$ .





$$\begin{split} \Phi_2(X,Y) &= X^3 + Y^3 - X^2 Y^2 + 1488 X^2 Y - 162000 X^2 + 1488 X Y^2 + 40773375 X Y \\ &\quad + 874800000 X - 162000 Y^2 + 8748000000 Y - 15746400000000 \end{split}$$



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# Mapping a volcano


Example  $\ell = 5$ , p = 4451, D = -151

General requirements  $4p = t^2 - v^2 \ell^2 D, \quad p \equiv 1 \mod \ell$ 



Example

 $\ell = 5, \quad p = 4451, \quad D = -151$  $t = 52, \quad v = 2, \quad h(D) = 7$   $\begin{array}{ll} \text{General requirements} \\ 4p = t^2 - v^2 \ell^2 D, \quad p \equiv 1 \bmod \ell \\ \ell \nmid v, \quad (\frac{D}{\ell}) = 1, \quad h(D) \geq \ell + 2 \end{array}$ 



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1. Find a root of  $H_D(X)$ 

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1. Find a root of  $H_D(X)$ : 901

Example  $\ell = 5$ , p = 4451, D = -151 t = 52, v = 2, h(D) = 7 $\ell_0 = 2$   $\begin{array}{l} \text{General requirements} \\ 4p = t^2 - v^2 \ell^2 D, \quad p \equiv 1 \mod \ell \\ \ell \nmid v, \quad (\frac{D}{\ell}) = 1, \quad h(D) \geq \ell + 2 \\ \ell_0 \neq \ell, \left(\frac{D}{\ell_0}\right) = 1 \end{array}$ 



2. Enumerate surface using the action of  $\alpha_{\ell_0}$ 

Example  $\ell = 5, \quad p = 4451, \quad D = -151$   $t = 52, \quad v = 2, \quad h(D) = 7$  $\ell_0 = 2, \quad \alpha_5 = \alpha_2^3$   $\begin{array}{l} \text{General requirements} \\ 4\rho = t^2 - v^2 \ell^2 D, \quad p \equiv 1 \ \text{mod} \ \ell \\ \ell \nmid v, \quad (\frac{D}{\ell}) = 1, \quad h(D) \geq \ell + 2 \\ \ell_0 \neq \ell, \ (\frac{D}{\ell_0}) = 1, \quad \alpha_\ell = \alpha_{\ell_0}^k \end{array}$ 



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3. Descend to the floor using Vélu's formula

 $\begin{array}{lll} \text{Example} & & \text{General re} \\ \ell &= 5, \quad p = 4451, \quad D = -151 & & 4\rho = t^2 \\ t &= 52, \quad v = 2, \quad h(D) = 7 & & \ell \nmid v, \quad (\ell_0 = 2, \quad \alpha_5 = \alpha_2^3 & & \ell_0 \neq \ell, \end{array}$ 





3. Descend to the floor using Vélu's formula: 901  $\stackrel{5}{\longrightarrow}$  3188

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4. Enumerate floor using the action of  $\beta_{\ell_0}$ 

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4. Enumerate floor using the action of  $\beta_{\ell_0}$  $\stackrel{2}{\longrightarrow}$  945  $\stackrel{2}{\longrightarrow}$  3144  $\stackrel{2}{\longrightarrow}$  3508  $\stackrel{2}{\longrightarrow}$  2843  $\stackrel{2}{\longrightarrow}$  1502  $\stackrel{2}{\longrightarrow}$  676  $\stackrel{2}{\longrightarrow}$  $\stackrel{2}{\longrightarrow}$  3497  $\stackrel{2}{\longrightarrow}$  1180  $\stackrel{2}{\longrightarrow}$  2464  $\stackrel{2}{\longrightarrow}$  4221  $\stackrel{2}{\longrightarrow}$  4228  $\stackrel{2}{\longrightarrow}$  2434  $\stackrel{2}{\longrightarrow}$  $\stackrel{2}{\longrightarrow}$  3244  $\stackrel{2}{\longrightarrow}$  2255  $\stackrel{2}{\longrightarrow}$  2976  $\stackrel{2}{\longrightarrow}$  3345  $\stackrel{2}{\longrightarrow}$  1064  $\stackrel{2}{\longrightarrow}$  1868  $\stackrel{2}{\longrightarrow}$  $\stackrel{2}{\longrightarrow}$  291  $\stackrel{2}{\longrightarrow}$  3147  $\stackrel{2}{\longrightarrow}$  2566  $\stackrel{2}{\longrightarrow}$  4397  $\stackrel{2}{\longrightarrow}$  2087  $\stackrel{2}{\longrightarrow}$  3341  $\stackrel{2}{\longrightarrow}$ 

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Example  $\ell = 5, \quad p = 4451, \quad D = -151$   $t = 52, \quad v = 2, \quad h(D) = 7$  $\ell_0 = 2, \quad \alpha_5 = \alpha_2^3, \quad \beta_{25} = \beta_2^7$ 





#### Interpolation



$$\begin{split} & \Phi_5(X, \ 901) = (X - \ 701)(X - \ 351)(X - \ 3188)(X - \ 2970)(X - \ 1478)(X - \ 3328) \\ & \Phi_5(X, \ 351) = (X - \ 901)(X - \ 2215)(X - \ 3508)(X - \ 2464)(X - \ 2976)(X - \ 2566) \\ & \Phi_5(X, \ 2215) = (X - \ 351)(X - \ 2501)(X - \ 3341)(X - \ 1868)(X - \ 2434)(X - \ 676) \\ & \Phi_5(X, \ 2501) = (X - \ 2215)(X - \ 2872)(X - \ 3147)(X - \ 2255)(X - \ 1180)(X - \ 3144) \\ & \Phi_5(X, \ 2872) = (X - \ 2501)(X - \ 1582)(X - \ 1502)(X - \ 4228)(X - \ 1064)(X - \ 2087) \\ & \Phi_5(X, \ 1582) = (X - \ 2872)(X - \ 701)(X - \ 945)(X - \ 3497)(X - \ 3244)(X - \ 291) \\ & \Phi_5(X, \ 701) = (X - \ 1582)(X - \ 901)(X - \ 2843)(X - \ 4221)(X - \ 3345)(X - \ 4397) \\ \end{split}$$

#### Interpolation



$$\begin{split} & \Phi_5(X, \ 901) = X^6 + 1337X^5 + \ 543X^4 + \ 497X^3 + 4391X^2 + 3144X + 3262 \\ & \Phi_5(X, \ 351) = X^6 + 3174X^5 + 1789X^4 + 3373X^3 + 3972X^2 + 2932X + 4019 \\ & \Phi_5(X, 2215) = X^6 + 2182X^5 + \ 512X^4 + \ 435X^3 + 2844X^2 + 2084X + 2709 \\ & \Phi_5(X, 2501) = X^6 + 2991X^5 + 3075X^5 + 3918X^3 + 2241X^2 + 3755X + 1157 \\ & \Phi_5(X, 2872) = X^6 + \ 389X^5 + 3292X^4 + 3909X^3 + \ 161X^2 + 1003X + 2091 \\ & \Phi_5(X, 1582) = X^6 + \ 1803X^5 + \ 794X^4 + \ 3584X^3 + \ 225X^2 + 1530X + 1975 \\ & \Phi_5(X, \ 701) = X^6 + \ 515X^5 + 1419X^4 + \ 941X^3 + 4145X^2 + 2722X + 2754 \end{split}$$

#### Interpolation



# Computing $\Phi_{\ell}(X, Y) \mod p$

Assume *D* and *p* are suitably chosen with  $D = O(\ell^2)$  and  $\log p = O(\log \ell)$ , and that  $H_D(X)$  has been precomputed.

- 1. Find a root of  $H_D(X)$  over  $\mathbb{F}_p$ .
- 2. Enumerate the surface(s) using cl(D)-action.
- 3. Descend to the floor using Vélu.
- 4. Enumerate the floor using  $cl(\ell^2 D)$ -action.
- 5. Build each  $\Phi_{\ell}(X, j_i)$  from its roots.
- 6. Interpolate  $\Phi_{\ell}(X, Y) \mod p$ .

Time complexity is  $O(\ell^2 \log^{3+\epsilon} \ell)$ . Space complexity is  $O(\ell^2 \log \ell)$ .  $\begin{array}{l} O(\ell \log^{3+\epsilon} \ell) \\ O(\ell \log^{2+\epsilon} \ell) \\ O(\ell \log^{1+\epsilon} \ell) \\ O(\ell^2 \log^{2+\epsilon} \ell) \\ O(\ell^2 \log^{3+\epsilon} \ell) \\ O(\ell^2 \log^{3+\epsilon} \ell) \\ O(\ell^2 \log^{3+\epsilon} \ell) \end{array}$ 

After computing  $\Phi_5(X, Y) \mod p$  for the primes:

4451, 6911, 9551, 28111, 54851, 110051, 123491, 160591, 211711, 280451, 434111, 530851, 686051, 736511, we apply the CRT to obtain

$$\begin{split} \Phi_5(X,Y) &= X^6 + Y^6 - X^5 Y^5 + 3720 (X^5 Y^4 + X^4 Y^5) - 4550940 (X^5 Y^3 + X^3 Y^5) \\ &+ 2028551200 (X^5 Y^2 + X^2 Y^5) - 246683410950 (X^5 Y + XY^5) + 1963211489280 (X^5 + Y^5) \\ &+ 1665999364600 X^4 Y^4 + 107878928185336800 (X^4 Y^3 + X^3 Y^4) \\ &+ 383083609779811215375 (X^4 Y^2 + X^2 Y^4) + 128541798906828816384000 (X^4 Y + XY^4) \\ &+ 1284733132841424456253440 (X^4 + Y^4) - 4550940 (X^3 Y^5 + X^5 Y^3) \\ &- 441206965512914835246100 X^3 Y^3 + 26898488858380731577417728000 (X^3 Y^2 + X^2 Y^3) \\ &- 192457934618928299655108231168000 (X^3 Y + XY^3) \\ &+ 280244777828439527804321565297868800 (X^3 + Y^3) \\ &+ 5110941777552418083110765199360000 X^2 Y^2 \\ &+ 36554736583949629295706472332656640000 (X^2 Y + XY^2) \\ &+ 6692500042627997708487149415015068467200 (X^2 + Y^2) \\ &- 264073457076620596259715790247978782949376XY \\ &+ 53274330803424425450420160273336509151232000 (X + Y) \\ &+ 141359947154721358697753474691071362751004672000. \end{split}$$

After computing  $\Phi_5(X, Y) \mod p$  for the primes:

4451, 6911, 9551, 28111, 54851, 110051, 123491, 160591, 211711, 280451, 434111, 530851, 686051, 736511, we apply the CRT to obtain

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(but note that  $\Phi_5^{f}(X, Y) = X^6 + Y^6 - X^5 Y^5 + 4XY$ ).

#### Computing $\Phi_\ell \mod m$

Given a prime  $\ell > 2$  and an integer m > 0:

- 1. Pick a discriminant D suitable for  $\ell$ .
- 2. Select a set of primes S suitable for  $\ell$  and D.
- 3. Precompute  $H_D$ , cl(D),  $cl(\ell^2 D)$ , and CRT data.
- 4. For each  $p \in S$ , compute  $\Phi_{\ell} \mod p$  and update CRT data.
- 5. Perform CRT postcomputation and output  $\Phi_{\ell} \mod m$ .

To compute  $\Phi_{\ell}$  over  $\mathbb{Z}$ , just use  $m = \prod p$ .

For "small" m, use explicit CRT mod m. For "large" m, standard CRT for large m. For m in between, use a hybrid approach.

# Complexity

#### Theorem (GRH)

For every prime  $\ell > 2$  there is a suitable discriminant D with  $|D| = O(\ell^2)$  for which there are  $\Omega(\ell^3 \log^3 \ell)$  primes  $p = O(\ell^6 (\log \ell)^4)$  that are suitable for  $\ell$  and D.

Heuristically,  $p = O(\ell^4)$ . In practice,  $\lg p < 64$ .

#### Theorem (GRH)

The expected running time is  $O(\ell^3 \log^3 \ell \log \log \ell)$ . The space required is  $O(\ell^2 \log(\ell m))$ .

#### An explicit height bound for $\Phi_\ell$

Let  $\ell$  be a prime. Let  $h(\Phi_{\ell})$  be the (natural) logarithmic height of  $\Phi_{\ell}$ .

Asymptotic bound:  $h(\Phi_{\ell}) = 6\ell \log \ell + O(\ell)$  (Paula Cohen 1984).

#### An explicit height bound for $\Phi_\ell$

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Asymptotic bound:  $h(\Phi_{\ell}) = 6\ell \log \ell + O(\ell)$  (Paula Cohen 1984).

Explicit bound:  $h(\Phi_{\ell}) \leq 6\ell \log \ell + 17\ell$  (Bröker-S 2009).

Conjectural bound:  $h(\Phi_{\ell}) \leq 6\ell \log \ell + 12\ell$  (for  $\ell > 30$ ).

The explicit bound holds for all  $\ell$ . The conjectural bound is known to hold for  $30 < \ell < 3600$ .

#### Other modular functions

We can compute polynomials relating f(z) and  $f(\ell z)$  for other modular functions, including the Weber f-function.

The coefficients of  $\Phi_{\ell}^{f}$  are roughly 72 times smaller. This means we need 72 fewer primes.

The polynomial  $\Phi_{\ell}^{f}$  is roughly 24 times sparser. This means we need 24 times fewer interpolation points.

We get a better than 1728-fold speedup using  $\Phi_{\ell}^{\dagger}$ .

#### Modular polynomials for $\ell = 11$

Classical:

 $X^{12} + Y^{12} - X^{11}Y^{11} + -1X^{11}Y^{11} + 8184X^{11}Y^{10} - 28278756X^{11}Y^9 + 53686822816X^{11}Y^8 + 28278756X^{11}Y^8 + 28278756X^{11}Y^9 + 282787575X^{11}Y^9 + 28277575X^{11}Y^{11}Y^{11}Y^{11}Y^{11}Y^{11}Y^{11}Y^{11}Y^{11}Y^{11}Y^{11}Y^{11}Y^{11}Y^{11}Y^{11}Y^$ 

- $\ 61058988656490 X^{11} Y^7 + 42570393135641712 X^{11} Y^6 17899526272883039048 X^{11} Y^5$
- $+\,4297837238774928467520 X^{11} Y^4\,-\,529134841844639613861795 X^{11} Y^3\,+\,27209811658056645815522600 X^{11} Y^2$
- $-\ 374642006356701393515817612 X^{11} Y + 296470902355240575283200000 X^{11}$
- ... 8 pages omitted ...
- + 392423345094527654908696 . . . 100 digits omitted . . . 000

#### Atkin:

$$\begin{split} X^{12} &- X^{11} Y + 744 X^{11} + 196680 X^{10} + 187 X^9 Y + 21354080 X^9 + 506 X^8 Y + 830467440 X^8 \\ &- 11440 X^7 Y + 16875327744 X^7 - 57442 X^6 Y + 208564958976 X^6 + 184184 X^5 Y + 1678582287360 X^5 \\ &+ 1675784 X^4 Y + 9031525113600 X^4 + 1867712 X^3 Y + 32349979904000 X^3 - 8252640 X^2 Y + 74246810880000 X^2 \\ &- 19849600 X Y + 98997734400000 X + Y^2 - 872000 Y + 58411072000000 \end{split}$$

#### Weber:

 $X^{12} + Y^{12} - X^{11}Y^{11} + 11X^9Y^9 - 44X^7Y^7 + 88X^5Y^5 - 88X^3Y^3 + 32XY^5$
## Weber modular polynomials

For  $\ell = 1009$ , the size of  $\Phi_{\ell}^{f}$  is 2.3MB, versus 3.9GB for  $\Phi_{\ell}$ , and computing  $\Phi_{\ell}^{f}$  takes 1.5s, versus 2840s for  $\Phi_{\ell}$ .

The current record is  $\ell = 60013$ . Working mod *m*, level  $\ell > 100000$  is feasible.

The polynomials  $\Phi^{f}_{\ell}$  for all  $\ell < 5000$  are available for download:

http://math.mit.edu/~drew