## Isogeny volcanoes: a computational perspective

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## A volcano



# A volcano



#### *l*-volcanoes

For a prime  $\ell$ , an  $\ell$ -volcano is a connected undirected graph whose vertices are partitioned into levels  $V_0, \ldots, V_d$ .

- **1.** The subgraph on  $V_0$  (the *surface*) is a connected regular graph of degree 0, 1, or 2.
- **2.** For i > 0, each  $v \in V_i$  has exactly one neighbor in  $V_{i-1}$ . All edges not on the surface arise in this manner.
- **3.** For i < d, each  $v \in V_i$  has degree  $\ell$ +1.

We allow self-loops and multi-edges in our graphs, but this can happen only on the surface of an  $\ell$ -volcano.

## A 3-volcano of depth 2



### **Elliptic curves**

An elliptic curve E/k is a smooth projective curve of genus 1 with a distinguished *k*-rational point 0.

For any field extension k'/k, the set of k'-rational points E(k') forms an abelian group with identity element 0.

When the characteristic of k is not 2 or 3 (which we assume for convenience) we may define E with an equation of the form

$$y^2 = x^3 + Ax + B,$$

where  $A, B \in k$ .

### *j*-invariants

The  $\bar{k}$ -isomorphism classes of elliptic curves E/k are in bijection with the field k. For  $E: y^2 = x^3 + Ax + B$ , the *j*-invariant of E is

$$j(E) = j(A, B) = 1728 \frac{4A^3}{4A^3 + 27B^2} \in k$$

The *j*-invariants j(0, B) = 0 and j(A, 0) = 1728 are special. They correspond to elliptic curves with extra automorphisms.

For  $j_0 \notin \{0, 1728\}$ , we have  $j_0 = j(A, B)$ , where

 $A = 3j_0(1728 - j_0)$  and  $B = 2j_0(1728 - j_0)^2$ .

Note that  $j(E_1) = j(E_2)$  does not necessarily imply that  $E_1$  and  $E_2$  are isomorphic over k, but they must be isomorphic over  $\bar{k}$ .

## $\ell$ -isogenies

An *isogeny*  $\phi: E_1 \rightarrow E_2$  is a morphism of elliptic curves, a rational map that fixes the point 0.

It induces a group homomorphism  $\phi: E_1(\bar{k}) \to E_2(\bar{k})$ . If  $\phi$  is nonzero then it has a finite kernel. Every finite subgroup of  $E_1(\bar{k})$  is the kernel of an isogeny.

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The *degree* of an isogeny is its degree as a rational map. For a nonzero *separable* isogeny,  $\deg \phi = |\ker \phi|$ .

We are interested in isogenies of prime degree  $\ell \neq \operatorname{char} k$ , which are necessarily separable isogenies with cyclic kernels.

The *dual isogeny*  $\hat{\phi} \colon E_2 \to E_1$  has the same degree  $\ell$  as  $\phi$ , and

$$\phi \circ \hat{\phi} = \hat{\phi} \circ \phi = [\ell]$$

is the *multiplication-by-* $\ell$  map.

## The $\ell$ -torsion subgroup

For  $\ell \neq \operatorname{char}(k)$ , the  $\ell$ -torsion subgroup

$$E[\ell] = \{P \in E(\bar{k}) : \ell P = 0\}$$

is isomorphic to  $\mathbb{Z}/\ell\mathbb{Z} \times \mathbb{Z}/\ell\mathbb{Z}$  and thus contains  $\ell + 1$  cyclic subgroups of order  $\ell$ , each of which is the kernel of an  $\ell$ -isogeny.

These  $\ell$ -isogenies are not necessarily defined over *k*.

An  $\ell$ -isogeny is defined over k (and has image defined over k) if and only if its kernel is Galois-invariant.

The number of Galois-invariant order- $\ell$  subgroups of  $E[\ell]$  is either 0, 1, 2, or  $\ell + 1$ .

### The modular equation

Let  $j: \mathbb{H} \to \mathbb{C}$  be the classical modular function. For any  $\tau \in \mathbb{H}$ , the values  $j(\tau)$  and  $j(\ell \tau)$  are the *j*-invariants of elliptic curves over  $\mathbb{C}$  that are  $\ell$ -isogenous.

The minimal polynomial  $\Phi_{\ell}(Y)$  of the function  $j(\ell z)$  over  $\mathbb{C}(j)$  has coefficients that are actually integer polynomials of j(z).

Replacing j(z) with X yields the *modular polynomial*  $\Phi_{\ell} \in \mathbb{Z}[X, Y]$  that parameterizes pairs of  $\ell$ -isogenous elliptic curves  $E/\mathbb{C}$ :

 $\Phi_{\ell}(j(E_1), j(E_2)) = 0 \iff j(E_1) \text{ and } j(E_2) \text{ are } \ell\text{-isogenous.}$ 

This moduli interpretation remains valid over any field of characteristic not  $\ell.$ 

 $<sup>\</sup>Phi_{\ell}(X,Y) = 0$  is a defining equation for the affine modular curve  $Y_0(\ell) = \Gamma_0(\ell) \setminus \mathbb{H}$ .

# The graph of $\ell$ -isogenies

#### Definition

The  $\ell$ -isogeny graph  $G_{\ell}(k)$  has vertex set  $\{j(E) : E/k\} = k$ and edges  $(j_1, j_2)$  for each root  $j_2 \in k$  of  $\Phi_{\ell}(j_1, Y)$  (with multiplicity).

Except for  $j \in \{0, 1728\}$ , the in-degree of each vertex of  $G_{\ell}$  is equal to its out-degree. Thus  $G_{\ell}$  is a bi-directed graph on  $k \setminus \{0, 1728\}$ , which we may regard as an undirected graph.

Note that we have an infinite family of graphs  $G_{\ell}(k)$  with vertex set k, one for each prime  $\ell \neq char(k)$ .

## Ordinary and supersingular curves

For an elliptic curve E/k with char(k) = p we have

 $E[p] \simeq \begin{cases} \mathbb{Z}/p\mathbb{Z} & (ordinary), \\ \{0\} & (supersingular). \end{cases}$ 

For isogenous elliptic curves  $E_1 \sim E_2$ , either both are ordinary or both are supersingular. Thus the each isogeny graph  $G_\ell$  decomposes into ordinary and supersingular components.

This has cryptographic applications; see [Charles-Lauter-Goren 2008], for example.

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Every supersingular curve is defined over  $\mathbb{F}_{p^2}$ . Thus the supersingular components of  $G_{\ell}(\mathbb{F}_{p^2})$  are regular graphs of degree  $\ell + 1$ .

In fact,  $G_{\ell}(\mathbb{F}_{p^2})$  has just one supersingular component, and it is a *Ramanujan graph* [Pizer 1990].

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## **Endomorphism rings**

Isogenies from an elliptic curve E to itself are *endomorphisms*. They form a ring End(E) under composition and point addition.

We always have  $\mathbb{Z} \subseteq \text{End}(E)$ , due to scalar multiplication maps. If  $\mathbb{Z} \subsetneq \text{End}(E)$ , then *E* has *complex multiplication* (CM).

For an elliptic curve *E* with complex multiplication:

 $\operatorname{End}(E) \simeq \begin{cases} \text{order in an imaginary quadratic field} & (\text{ordinary}), \\ \text{order in a quaternion algebra} & (\text{supersingular}). \end{cases}$ 

Every elliptic curve over a finite field  $\mathbb{F}_q$  has CM, since if *E* is ordinary then the *Frobenius endomorphism*  $\pi_E(x, y) = (x^q, y^q)$  does not lie in  $\mathbb{Z}$ .

#### Horizontal and vertical isogenies

Let  $\varphi \colon E_1 \to E_2$  by an  $\ell$ -isogeny of ordinary elliptic curves with CM. Let  $\operatorname{End}(E_1) \simeq \mathcal{O}_1 = [1, \tau_1]$  and  $\operatorname{End}(E_2) \simeq \mathcal{O}_2 = [1, \tau_2]$ .

Then  $\ell \tau_2 \in \mathcal{O}_1$  and  $\ell \tau_1 \in \mathcal{O}_2$ .

Thus one of the following holds:

- $\mathcal{O}_1 = \mathcal{O}_2$ , in which case  $\varphi$  is *horizontal*;
- $[\mathcal{O}_1 : \mathcal{O}_2] = \ell$ , in which case  $\varphi$  is *descending*;
- $[\mathcal{O}_2 : \mathcal{O}_1] = \ell$ , in which case  $\varphi$  is *ascending*.

In the latter two cases we say that  $\varphi$  is a *vertical* isogeny.

## The theory of complex multiplication

Let E/k have  $\operatorname{End}(E) \simeq \mathcal{O} \subset K = \mathbb{Q}(\sqrt{D})$ , with  $D = \operatorname{disc} K$ .

For each invertible  $\mathcal{O}\text{-ideal}\ \mathfrak{a},$  the  $\mathfrak{a}\text{-torsion}\ subgroup$ 

$$E[\mathfrak{a}] = \{ P \in E(\bar{k}) : \alpha(P) = 0 \text{ for all } \alpha \in \mathfrak{a} \}$$

is the kernel of an isogeny  $\varphi_{\mathfrak{a}} : E \to E'$  of degree  $N(\mathfrak{a}) = [\mathcal{O} : \mathfrak{a}]$ . We necessarily have  $\operatorname{End}(E) \simeq \operatorname{End}(E')$ , so  $\varphi_{\mathfrak{a}}$  is **horizontal**.

If a is principal, then  $E' \simeq E$ . This induces a  $cl(\mathcal{O})$ -action on the set.

$$\operatorname{Ell}_{\mathcal{O}}(k) = \{j(E) : E/k \text{ with } \operatorname{End}(E) \simeq \mathcal{O}\}.$$

This action is faithful and transitive; thus  $\text{Ell}_{\mathcal{O}}(k)$  is a principal homogeneous space, a *torsor*, for  $\text{cl}(\mathcal{O})$ .

One can decompose horizontal isogenies of large prime degree into an equivalent sequence of isogenies of small prime degrees, which makes them **easy to compute**; see [Bröker-Charles-Lauter 2008, Jao-Souhkarev ANTS IX].

### **Horizontal isogenies**

Every horizontal  $\ell\text{-isogeny}$  arises from the action of an invertible  $\mathcal O\text{-ideal }\mathfrak l$  of norm  $\ell.$ 

If  $\ell \mid [\mathcal{O}_K : \mathcal{O}]$ , no such  $\mathfrak{l}$  exists; if  $\ell \nmid [\mathcal{O}_K : \mathcal{O}]$ , then there are

$$1 + \left(\frac{D}{\ell}\right) = \begin{cases} 0\\1\\2 \end{cases}$$

 $\ell$  is inert in *K*,  $\ell$  is ramified in *K*,  $\ell$  splits in *K*,

such  $\ell$ -isogenies.

In the split case,  $(\ell) = \mathfrak{l} \cdot \overline{\mathfrak{l}}$ , and the  $\mathfrak{l}$ -orbits partition  $\operatorname{Ell}_{\mathcal{O}}(k)$  into cycles corresponding to the cosets of  $\langle [\mathfrak{l}] \rangle$  in  $\operatorname{cl}(\mathcal{O})$ .

#### **Vertical isogenies**

Let  $\mathcal{O}$  be an imaginary quadratic order with discriminant  $D_{\mathcal{O}} < -4$ , and let  $\mathcal{O}' = \mathbb{Z} + \ell \mathcal{O}$  be the order of index  $\ell$  in  $\mathcal{O}$ .

The map that sends each invertible  $\mathcal{O}'$ -ideal  $\mathfrak{a}$  to the (invertible)  $\mathcal{O}$ -ideal  $\mathfrak{a}\mathcal{O}$  preserves norms and induces a surjective homomorphism

 $\phi\colon \operatorname{cl}(\mathcal{O}')\to\operatorname{cl}(\mathcal{O})$ 

compatible with the class group actions on  $\text{Ell}_{\mathcal{O}}(k)$  and  $\text{Ell}_{\mathcal{O}'}(k)$ .

It follows that each  $j(E') \in \text{Ell}_{\mathcal{O}'}(k)$  has a unique  $\ell$ -isogenous "parent" j(E) in  $\text{Ell}_{\mathcal{O}}(k)$ , and every vertical isogeny must arise in this way.

The "children" of j(E) correspond to a coset of the kernel of  $\phi$ , which is a cyclic of order  $\ell - \left(\frac{D_{\mathcal{O}}}{\ell}\right)$ , generated by the class of an invertible  $\mathcal{O}'$ -ideal with norm  $\ell^2$ .

#### Ordinary elliptic curves over finite fields

Let  $E/\mathbb{F}_q$  be an ordinary elliptic curve with *trace of Frobenius* 

 $t = \operatorname{tr} \pi_E = q + 1 - \# E(\mathbb{F}_q).$ 

Then  $\pi_E^2 - t\pi_E + q = 0$  and we have the *norm equation* 

$$4q = t^2 - v^2 D,$$

where *D* is the (fundamental) discriminant of the imaginary quadratic field  $K = \mathbb{Q}(\sqrt{t^2 - 4q}) \simeq \operatorname{End}(E) \otimes \mathbb{Q}$  and  $v = [\mathcal{O}_K : \mathbb{Z}[\pi_E]]$ . We have

 $\mathbb{Z}[\pi_E] \subseteq \operatorname{End}(E) \subseteq \mathcal{O}_K.$ 

Thus  $[\mathcal{O}_K : \text{End}(E)]$  divides *v*; this holds for any *E* with trace *t*. If we define  $\text{Ell}_t(\mathbb{F}_q) = \{j(E) : E/\mathbb{F}_q \text{ with } \text{tr } \pi_E = t\}$ , then

$$\operatorname{Ell}_{t}(\mathbb{F}_{q}) = \bigcup_{\mathbb{Z}[\pi_{E}] \subseteq \mathcal{O} \subseteq \mathcal{O}_{K}} \operatorname{Ell}_{\mathcal{O}}(\mathbb{F}_{q}).$$

### The main theorem

#### **Theorem (Kohel)**

Let *V* be an ordinary connected component of  $G_{\ell}(\mathbb{F}_q)$  that does not contain 0, 1728. Then *V* is an  $\ell$ -volcano in which the following hold:

(i) Vertices in level  $V_i$  all have the same endomorphism ring  $\mathcal{O}_i$ .

(ii) 
$$\ell \nmid [\mathcal{O}_K : \mathcal{O}_0]$$
, and  $[\mathcal{O}_i : \mathcal{O}_{i+1}] = \ell$ .

- (iii) The subgraph on  $V_0$  has degree  $1 + (\frac{D}{\ell})$ , where  $D = \operatorname{disc}(\mathcal{O}_0)$ .
- (iv) If  $(\frac{D}{\ell}) \ge 0$  then  $|V_0|$  is the order of [I] in  $cl(\mathcal{O}_0)$ .
- (v) The depth of V is  $\operatorname{ord}_{\ell}(v)$ , where  $4q = t^2 v^2 D$ .

The term volcano is due to Fouquet and Morain (ANTS V).

# **Applications**













Curves on the floor necessarily have cyclic rational  $\ell$ -torsion. This is useful, for example, when constructing Edwards curves with the CM method [Morain 2009].

#### Finding a shortest path to the floor



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We now know that we are 2 levels above the floor.

# Application: identifying supersingular curves

The equation  $4q = t^2 - v^2D$  implies that each ordinary component of  $G_{\ell}(\mathbb{F}_q)$  is an  $\ell$ -volcano of depth less than  $\log_{\ell} \sqrt{4q}$ .

Given  $j(E) \in \mathbb{F}_{p^2}$ , if we cannot find a shortest path to the floor in  $G_2(\mathbb{F}_{p^2})$  within  $\lceil \log_2 p \rceil$  steps, then *E* **must be supersingular**.

Conversely, if *E* is supersingular, our attempt to find the floor must fail, since every vertex in the supersingular component has degree  $\ell + 1$ .

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Conversely, if *E* is supersingular, our attempt to find the floor must fail, since every vertex in the supersingular component has degree  $\ell + 1$ .

This yields a (probabilistic) algorithm to determine supersingularity in  $\tilde{O}(n^3)$  time, where  $n = \log p$ , improving the  $\tilde{O}(n^4)$  complexity of the best previously known algorithms.

Moreover, the expected running time on a random elliptic curve is  $\tilde{O}(n^2)$ , matching the complexity of the best *Monte Carlo* algorithms, and faster in practice.

See [S 2012] for details.

# Application: computing endomorphism rings

Given an ordinary elliptic curve  $E/\mathbb{F}_q$ , if we compute the Frobenius trace *t* and put  $4q = t^2 - v^2D$ , we can determine  $\mathcal{O} \simeq \text{End}(E)$  by determining  $u = [\mathcal{O}_K : \mathcal{O}]$ , which must divide *v*.

It suffices to determine the level of j(E) in its  $\ell$ -volcano for  $\ell | v$ .

**Problem**: when  $\ell$  is large it is not feasible to compute  $\Phi_{\ell}$ , nor is it feasible to directly compute a **vertical**  $\ell$ -isogeny.

See [Bisson-S 2011] and [Bisson 2011] for more details.

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**Solution**: we may determine the primes  $\ell | u$  by finding *smooth relations* that hold in  $cl((v/\ell)^2D)$  but not in  $cl(\ell^2D)$  and evaluating the corresponding **horizontal** isogenies (and similarly for  $\ell^e$ )

This yields a probabilistic algorithm to compute End(E) with subexponential expected running time  $L[1/2, \sqrt{3}/2]$ , under GRH.

See [Bisson-S 2011] and [Bisson 2011] for more details.

#### Example

Let  $q = 2^{320} + 261$  and suppose tr  $\pi_E = t$ , where t = 2306414344576213633891236434392671392737040459558.

Then  $4q = t^2 - v^2 D$ , where D = -147759 and  $v = 2^2 p_1 p_2$  with

 $p_1 = 16447689059735824784039,$  $p_2 = 71003976975490059472571.$ 

For  $D_1 = 2^4 p_2^2 D$ , and  $D'_1 = p_1^2 D$ , the relation

 $\{\mathfrak{p}_5,\mathfrak{p}_{19}^2,\bar{\mathfrak{p}}_{23}^{210},\mathfrak{p}_{29},\mathfrak{p}_{31},\bar{\mathfrak{p}}_{41}^{145},\mathfrak{p}_{139},\bar{\mathfrak{p}}_{149},\mathfrak{p}_{167},\bar{\mathfrak{p}}_{191},\bar{\mathfrak{p}}_{251}^6,\mathfrak{p}_{269},\bar{\mathfrak{p}}_{587}^7,\bar{\mathfrak{p}}_{643}\}$ 

holds in  $cl(D_1)$  but not in  $cl(D'_1)$  ( $\mathfrak{p}_{\ell}$  is an ideal of norm  $\ell$ ). For  $D_2 = 2^4 p_1^2 D$ , and  $D'_2 = p_2^2 D$ , the relation

 $\{\mathfrak{p}_{11}, \bar{\mathfrak{p}}_{13}^{576}, \mathfrak{p}_{23}^2, \bar{\mathfrak{p}}_{41}, \bar{\mathfrak{p}}_{47}, \mathfrak{p}_{83}, \mathfrak{p}_{101}, \bar{\mathfrak{p}}_{197}^{28}, \bar{\mathfrak{p}}_{307}^3, \mathfrak{p}_{317}, \bar{\mathfrak{p}}_{419}, \mathfrak{p}_{911}\}$ 

holds in  $cl(D_2)$  but not in  $cl(D'_2)$ .

### Constructing elliptic curves with the CM method

Let  $\mathcal{O}$  be an imaginary quadratic order with discriminant D. The *Hilbert class polynomial*  $H_D \in \mathbb{Z}[X]$  is defined by

$$H_D(X) = \prod_{j \in \operatorname{Ell}_{\mathcal{O}}(\mathbb{C})} (X - j).$$

Equivalently, it is the minimal polynomial of  $j(\mathcal{O})$  over  $K = \mathbb{Q}(\sqrt{D})$ . The field  $K_{\mathcal{O}} = K(j(\mathcal{O}))$  is the *ring class field* for  $\mathcal{O}$ .

One can also construct supersingular curves with Hilbert class polynomials; see [Bröker 2008].
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If *q* splits completely in  $K_{\mathcal{O}}$ , then  $H_D(X)$  splits completely in  $\mathbb{F}_q[X]$ , and every root of  $H_D$  is the *j*-invariant of an elliptic curve  $E/\mathbb{F}_q$  with N = q + 1 - t points, where  $4q = t^2 - v^2D$ .

Every ordinary elliptic curve  $E/\mathbb{F}_q$  can be constructed in this way, but computing  $H_D$  becomes quite difficult as |D| grows.

The size of  $H_D$  is  $O(|D| \log |D|)$  bits, exponential in  $\log q$ .

One can also construct supersingular curves with Hilbert class polynomials; see [Bröker 2008].

### Application: computing Hilbert class polynomials

The CRT approach to computing  $H_D(X)$ , as described in [Belding-Bröker-Enge-Lauter ANTX VIII] and [S 2011].

- **1.** Select a sufficiently large set of primes of the form  $4p = t^2 v^2D$ .
- **2.** For each prime p, compute  $H_D \mod p$  as follows:
  - **a.** Generate random curves  $E/\mathbb{F}_p$  until #E = p + 1 t.
  - **b.** Use volcano climbing to find  $E' \sim E$  with  $\operatorname{End}(E') \simeq \mathcal{O}$ .
  - **c.** Enumerate  $\text{Ell}_{\mathcal{O}}(\mathbb{F}_p)$  by applying the  $\text{cl}(\mathcal{O})$ -action to j(E').
  - **d.** Compute  $\prod_{j \in \text{Ell}_{\mathcal{O}}(\mathbb{F}_p)} (X j) = H_D(X) \mod p$ .
- **3.** Use the CRT to recover  $H_D$  over  $\mathbb{Z}$  (or mod q via the explicit CRT).

Under the GRH, the expected running time is  $O(|D| \log^{5+\epsilon} |D|)$ , quasi-linear in the size of  $H_D$ .

One can similarly compute other types of class polynomials [Enge-S ANTS IX].

## Using a polycyclic presentation

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But cl(D) is also generated by the classes of ideals  $\mathfrak{a}_2$  and  $\mathfrak{a}_{13}$  with norms 2 and 13. The classes  $[\mathfrak{a}_2]$  and  $[\mathfrak{a}_{13}]$  have orders 20 and 50 and thus are not independent in  $cl(\mathcal{O})$ , in fact  $[\mathfrak{a}_{13}]^5 = [\mathfrak{a}_2]^{18}$ .

Nevertheless, every element of cl(D) can uniquely represented as

 $[\mathfrak{a}_2]^{e_2}[\mathfrak{a}_{13}]^{e_{13}},$ 

with  $0 \le e_2 < 20$  and  $0 \le e_{13} < 5$ .

In general, any sequence of generators for a finite abelian group *G* determines a *polycyclic presentation* for *G*.

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In general, any sequence of generators for a finite abelian group *G* determines a *polycyclic presentation* for *G*.

Using the polycyclic presentation  $([\mathfrak{a}_2], [\mathfrak{a}_{13}])$  to enumerate  $\text{Ell}_{\mathcal{O}}(\mathbb{F}_p)$  is **100 times faster** than using  $([\mathfrak{a}_{19}])$ .















For particularly deep volcanoes, one may prefer to use a pairing-based approach; see [Ionica-Joux ANTS IX].

#### **Computational results**

The CRT method has been used to compute  $H_D(X)$  with  $|D| > 10^{13}$ , and using alternative class polynomials, with  $|D| > 10^{15}$  (for comparison, the previous record was  $|D| \approx 10^{10}$ ).

When  $cl(\mathcal{O})$  is composite (almost always the case), one can accelerate the CM method by decomposing the ring class field [Hanrot-Morain 2001, Enge-Morain 2003].

Combining this idea with the CRT approach has made CM constructions with  $|D| > 10^{16}$  possible [S 2012].

#### Application: computing modular polynomials

We can also use a CRT approach to compute  $\Phi_{\ell}(X, Y)$  [Bröker-Lauter-S 2012].

- **1.** Select a sufficiently large set of primes of the form  $4p = t^2 \ell^2 v^2 D$  with  $\ell \nmid v, p \equiv 1 \mod \ell$ , and  $h(D) > \ell + 1$ .
- **2.** For each prime *p*, compute  $\Phi_{\ell} \mod p$  as follows:
  - **a.** Compute  $\operatorname{Ell}_{\mathcal{O}}(\mathbb{F}_p)$  using  $H_D \mod p$ .
  - **b.** Map the  $\ell$ -volcanoes intersecting  $\text{Ell}_O(\mathbb{F}_p)$  (without using  $\Phi_\ell$ ).
  - **c.** Interpolate  $\Phi_{\ell}(X, Y) \mod p$ .
- **3.** Use the CRT to recover  $\Phi_{\ell}$  over  $\mathbb{Z}$  (or mod q via the explicit CRT).

Under the GRH, the expected running time is  $O(\ell^3 \log^{3+\epsilon} \ell)$ , quasi-linear in the size of  $\Phi_{\ell}$ .

We can similarly compute modular polynomials for other modular functions. One can also use a CRT approach to compute  $\Phi_N$  for composite *N* [Ono-S in prog].



Example  $\ell = 5$ , p = 4451, D = -151

General requirements  $4p = t^2 - v^2 \ell^2 D, \quad p \equiv 1 \mod \ell$ 



Example  $\ell = 5, p = 4451, D = -151$ t = 52, v = 2, h(D) = 7  $\begin{array}{ll} \mbox{General requirements} \\ 4p = t^2 - v^2 \ell^2 D, \quad p \equiv 1 \mbox{ mod } \ell \\ \ell \nmid v, \quad \left( \frac{D}{\ell} \right) = 1, \quad h(D) \geq \ell + 2 \end{array}$ 



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1. Find a root of  $H_D(X)$ 

Example  $\ell = 5$ , p = 4451, D = -151t = 52, v = 2, h(D) = 7  $\begin{array}{ll} \mbox{General requirements} \\ 4p = t^2 - v^2 \ell^2 D, \quad p \equiv 1 \mbox{ mod } \ell \\ \ell \nmid v, \quad \left( \frac{D}{\ell} \right) = 1, \quad h(D) \geq \ell + 2 \end{array}$ 



1. Find a root of  $H_D(X)$ : 901





2. Enumerate surface using the action of  $\alpha_{\ell_0}$ 



























 $\begin{array}{l} \mbox{General requirements} \\ 4p = t^2 - v^2 \ell^2 D, \quad p \equiv 1 \mbox{ mod } \ell \\ \ell \nmid v, \quad (\frac{D}{\ell}) = 1, \quad h(D) \geq \ell + 2 \\ \ell_0 \neq \ell, \ (\frac{D}{\ell_0}) = 1, \ \alpha_\ell = \alpha_{\ell_0}^k \end{array}$ 









3. Descend to the floor using Vélu's formula





3. Descend to the floor using Vélu's formula: 901  $\stackrel{5}{\longrightarrow}$  3188



 $\begin{array}{l} \mbox{General requirements} \\ 4p = t^2 - \nu^2 \ell^2 D, \quad p \equiv 1 \mbox{ mod } \ell \\ \ell \nmid \nu, \quad (\frac{D}{\ell}) = 1, \quad h(D) \geq \ell + 2 \\ \ell_0 \neq \ell, \ (\frac{D}{\ell_0}) = 1, \ \alpha_\ell = \alpha_{\ell_0}^k \end{array}$ 



4. Enumerate floor using the action of  $\beta_{\ell_0}$ 





4. Enumerate floor using the action of  $\beta_{\ell_0}$   $3188 \xrightarrow{2}{2} 945 \xrightarrow{2}{2} 3144 \xrightarrow{2}{2} 3508 \xrightarrow{2}{2} 2843 \xrightarrow{2}{2} 1502 \xrightarrow{2}{2} 676 \xrightarrow{2}{2}$   $2970 \xrightarrow{2}{2} 3497 \xrightarrow{2}{2} 1180 \xrightarrow{2}{2} 2464 \xrightarrow{2}{2} 4221 \xrightarrow{2}{2} 4228 \xrightarrow{2}{2} 2434 \xrightarrow{2}{2}$   $1478 \xrightarrow{2}{2} 3244 \xrightarrow{2}{2} 2255 \xrightarrow{2}{2} 2976 \xrightarrow{2}{2} 3345 \xrightarrow{2}{2} 1064 \xrightarrow{2}{2} 1868 \xrightarrow{2}{2}$  $3328 \xrightarrow{2}{2} 291 \xrightarrow{2}{3147} \xrightarrow{2} 2566 \xrightarrow{2}{4397} \xrightarrow{2} 2087 \xrightarrow{2}{3341} \xrightarrow{2}$ 





4. Enumerate floor using the action of  $\beta_{\ell_0}$  $\xrightarrow{2}{2}$  945  $\xrightarrow{2}{2}$  3144  $\xrightarrow{2}{2}$  3508  $\xrightarrow{2}{2}$  2843  $\xrightarrow{2}{2}$  1502  $\xrightarrow{2}{2}$  676  $\xrightarrow{2}{2}$  $\xrightarrow{2}{2}$  3497  $\xrightarrow{2}{2}$  1180  $\xrightarrow{2}{2}$  2464  $\xrightarrow{2}{2}$  4221  $\xrightarrow{2}{2}$  4228  $\xrightarrow{2}{2}$  2434  $\xrightarrow{2}{2}$  $\xrightarrow{2}{2}$  3244  $\xrightarrow{2}{2}$  2255  $\xrightarrow{2}{2}$  2976  $\xrightarrow{2}{2}$  3345  $\xrightarrow{2}{2}$  1064  $\xrightarrow{2}{2}$  1868  $\xrightarrow{2}{2}$  $\xrightarrow{2}{2}$  291  $\xrightarrow{2}{3}$  3147  $\xrightarrow{2}{2}$  2566  $\xrightarrow{2}{3}$  4397  $\xrightarrow{2}{2}$  2087  $\xrightarrow{2}{3}$  3341  $\xrightarrow{2}{2}$ 





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4. Enumerate floor using the action of  $\beta_{\ell_0}$  $3188 \xrightarrow{2}{2} 945 \xrightarrow{2}{2} 3144 \xrightarrow{2}{2} 3508 \xrightarrow{2}{2} 2843 \xrightarrow{2}{2} 1502 \xrightarrow{2}{2} 676 \xrightarrow{2}{2} 2970 \xrightarrow{2}{2} 3497 \xrightarrow{2}{2} 1180 \xrightarrow{2}{2} 2464 \xrightarrow{2}{2} 4221 \xrightarrow{2}{2} 4228 \xrightarrow{2}{2} 2434 \xrightarrow{2}{2} 1478 \xrightarrow{2}{2} 3244 \xrightarrow{2}{2} 2255 \xrightarrow{2}{2} 2976 \xrightarrow{2}{2} 3345 \xrightarrow{2}{2} 1064 \xrightarrow{2}{2} 1868 \xrightarrow{2}{2} 3328 \xrightarrow{2}{2} 291 \xrightarrow{2}{3} 3147 \xrightarrow{2} 2566 \xrightarrow{2}{3} 4397 \xrightarrow{2}{2} 2087 \xrightarrow{3}{3} 341 \xrightarrow{2}$ 





4. Enumerate floor using the action of  $\beta_{\ell_0}$   $3188 \xrightarrow{2}{2} 945 \xrightarrow{2}{2} 3144 \xrightarrow{2}{2} 3508 \xrightarrow{2}{2} 2843 \xrightarrow{2}{2} 1502 \xrightarrow{2}{2} 676 \xrightarrow{2}{2}$   $2970 \xrightarrow{2}{2} 3497 \xrightarrow{2}{2} 1180 \xrightarrow{2}{2} 2464 \xrightarrow{2}{2} 4221 \xrightarrow{2}{2} 4228 \xrightarrow{2}{2} 2434 \xrightarrow{2}{2}$   $1478 \xrightarrow{2}{2} 3244 \xrightarrow{2}{2} 2255 \xrightarrow{2}{2} 2976 \xrightarrow{2}{2} 3345 \xrightarrow{2}{2} 1064 \xrightarrow{2}{2} 1868 \xrightarrow{2}{2}$  $3328 \xrightarrow{2}{2} 291 \xrightarrow{2}{3} 3147 \xrightarrow{2} 2566 \xrightarrow{2}{4} 4397 \xrightarrow{2}{2} 2087 \xrightarrow{2}{3} 3341 \xrightarrow{2}{2}$ 





4. Enumerate floor using the action of  $\beta_{\ell_0}$   $3188 \xrightarrow{2}{2} 945 \xrightarrow{2}{2} 3144 \xrightarrow{2}{2} 3508 \xrightarrow{2}{2} 2843 \xrightarrow{2}{2} 1502 \xrightarrow{2}{2} 676 \xrightarrow{2}{2}$   $2970 \xrightarrow{2}{2} 3497 \xrightarrow{2}{2} 1180 \xrightarrow{2}{2} 2464 \xrightarrow{2}{2} 4221 \xrightarrow{2}{2} 4228 \xrightarrow{2}{2} 2434 \xrightarrow{2}{2}$   $1478 \xrightarrow{2}{2} 3244 \xrightarrow{2}{2} 2255 \xrightarrow{2}{2} 2976 \xrightarrow{2}{2} 3345 \xrightarrow{2}{2} 1064 \xrightarrow{2}{2} 1868 \xrightarrow{2}{2}$  $3328 \xrightarrow{2}{2} 291 \xrightarrow{2}{3147} \xrightarrow{2} 2566 \xrightarrow{2}{4397} \xrightarrow{2} 2087 \xrightarrow{2}{3341} \xrightarrow{2}$
# Mapping a volcano





4. Enumerate floor using the action of  $\beta_{\ell_0}$  $3188 \xrightarrow{2}{2} 945 \xrightarrow{2}{2} 3144 \xrightarrow{2}{2} 3508 \xrightarrow{2}{2} 2843 \xrightarrow{2}{2} 1502 \xrightarrow{2}{2} 676 \xrightarrow{2}{2} 2970 \xrightarrow{2}{3497} \xrightarrow{2}{1180} \xrightarrow{2}{2} 2464 \xrightarrow{2}{2} 4221 \xrightarrow{2}{2} 4228 \xrightarrow{2}{2} 2434 \xrightarrow{2}{1478} \xrightarrow{2}{3244} \xrightarrow{2}{2} 2255 \xrightarrow{2}{2} 2976 \xrightarrow{2}{3345} \xrightarrow{2}{2} 1064 \xrightarrow{2}{2} 1868 \xrightarrow{2}{328} \xrightarrow{2}{291} \xrightarrow{3147} \xrightarrow{2} 2566 \xrightarrow{2}{3397} \xrightarrow{2} 2087 \xrightarrow{2}{3341} \xrightarrow{2}$ 

# Mapping a volcano





4. Enumerate floor using the action of  $\beta_{\ell_0}$   $3188 \xrightarrow{2}{2} 945 \xrightarrow{2}{2} 3144 \xrightarrow{2}{2} 3508 \xrightarrow{2}{2} 2843 \xrightarrow{2}{2} 1502 \xrightarrow{2}{2} 676 \xrightarrow{2}{2}$   $2970 \xrightarrow{2}{2} 3497 \xrightarrow{2}{2} 1180 \xrightarrow{2}{2} 2464 \xrightarrow{2}{2} 4221 \xrightarrow{2}{2} 4228 \xrightarrow{2}{2} 2434 \xrightarrow{2}{2}$   $1478 \xrightarrow{2}{2} 3244 \xrightarrow{2}{2} 2255 \xrightarrow{2}{2} 2976 \xrightarrow{2}{2} 3345 \xrightarrow{2}{2} 1064 \xrightarrow{2}{2} 1868 \xrightarrow{2}{2}$  $3328 \xrightarrow{2}{2} 291 \xrightarrow{2}{3147} \xrightarrow{2} 2566 \xrightarrow{4}{4397} \xrightarrow{2} 2087 \xrightarrow{2} 3341 \xrightarrow{2}$ 

# Mapping a volcano





### Interpolating $\Phi_\ell \mod p$



$$\begin{split} \Phi_5(X, \ 901) &= (X - \ 701)(X - \ 351)(X - \ 3188)(X - \ 2970)(X - \ 1478)(X - \ 3328) \\ \Phi_5(X, \ 351) &= (X - \ 901)(X - \ 2215)(X - \ 3508)(X - \ 2464)(X - \ 2976)(X - \ 2566) \\ \Phi_5(X, \ 2215) &= (X - \ 351)(X - \ 2501)(X - \ 3341)(X - \ 1868)(X - \ 2434)(X - \ 676) \\ \Phi_5(X, \ 2501) &= (X - \ 2215)(X - \ 2872)(X - \ 3147)(X - \ 2255)(X - \ 1180)(X - \ 3144) \\ \Phi_5(X, \ 2872) &= (X - \ 2501)(X - \ 1582)(X - \ 1502)(X - \ 4228)(X - \ 1064)(X - \ 2087) \\ \Phi_5(X, \ 1582) &= (X - \ 2872)(X - \ 701)(X - \ 945)(X - \ 3497)(X - \ 3244)(X - \ 291) \\ \Phi_5(X, \ 701) &= (X - \ 1582)(X - \ 901)(X - \ 2843)(X - \ 4221)(X - \ 3345)(X - \ 4397) \\ \end{split}$$

### Interpolating $\Phi_\ell \mod p$



$$\begin{split} \Phi_5(X, \ 901) &= X^6 + 1337X^5 + 543X^4 + 497X^3 + 4391X^2 + 3144X + 3262 \\ \Phi_5(X, \ 351) &= X^6 + 3174X^5 + 1789X^4 + 3373X^3 + 3972X^2 + 2932X + 4019 \\ \Phi_5(X, 2215) &= X^6 + 2182X^5 + 512X^4 + 435X^3 + 2844X^2 + 2084X + 2709 \\ \Phi_5(X, 2501) &= X^6 + 2991X^5 + 3075X^5 + 3918X^3 + 2241X^2 + 3755X + 1157 \\ \Phi_5(X, 2872) &= X^6 + 389X^5 + 3292X^4 + 3909X^3 + 161X^2 + 1003X + 2091 \\ \Phi_5(X, 1582) &= X^6 + 1803X^5 + 794X^4 + 3584X^3 + 225X^2 + 1530X + 1975 \\ \Phi_5(X, \ 701) &= X^6 + 515X^5 + 1419X^4 + 941X^3 + 4145X^2 + 2722X + 2754 \end{split}$$

# Interpolating $\Phi_\ell \mod p$



$$\begin{split} \Phi_5(X,Y) &= X^6 + (4450Y^5 + 3720Y^4 + 2433Y^3 + 3499Y^2 + & 70Y + 3927)X^5 \\ &(3720Y^5 + 3683Y^4 + 2348Y^3 + 2808Y^2 + 3745Y + & 233)X^4 \\ &(2433Y^5 + 2348Y^4 + 2028Y^3 + 2025Y^2 + 4006Y + 2211)X^3 \\ &(3499Y^5 + 2808Y^4 + 2025Y^3 + 4378Y^2 + 3886Y + 2050)X^2 \\ &(& 70Y^5 + 3745Y^4 + 4006Y^3 + 3886Y^2 + & 905Y + 2091)X \\ &(& Y^6 + & 3927Y^5 + & 233Y^4 + 2211Y^3 + 2050Y^2 + 2091Y + 2108) \end{split}$$

## **Computational results**

#### Level records

- **1. 10009**:  $\Phi_\ell$
- **2.** 20011:  $\Phi_{\ell} \mod q$
- **3.** 60013:  $\Phi_{\ell}^{f}$

#### Speed records

- **1. 251**:  $\Phi_{\ell}$  in 28s  $\Phi_{\ell}$  mod q in 4.8s (vs 688s)
- **2.** 1009:  $\Phi_{\ell}$  in 2830s  $\Phi_{\ell} \mod q$  in 265s (vs 107200s)
- **3.** 1009:  $\Phi_{\ell}^{f}$  in 2.8s

Effective throughput when computing  $\Phi_{1009} \mod q$  is 100Mb/s.

Single core CPU times (AMD 3.0 GHz), using prime  $q \approx 2^{256}$ . Polynomials  $\Phi_{\ell}^{f}$  for  $\ell < 5000$  available at http://math.mit.edu/~drew.